## Outline for Today.

Polynomials.

Secret Sharing.

Finite Fields.

Share secret among *n* people.

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Two points make a line. Lots of lines go through one point.

## Polynomials

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is specified by **coefficients**  $a_d, \dots a_0$ . P(x) **contains** point (a, b) if b = P(a). **Polynomials over reals**:  $a_1, \dots, a_d \in \Re$ , use  $x \in \Re$ .

E.g., Reals, rationals, complex numbers.

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Intuitively, a field is a set with operations corresponding to addition, multiplication, and division.



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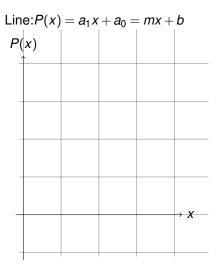
Polynomials P(x) with arithmetic modulo p:

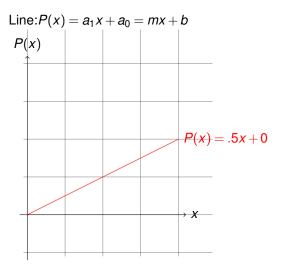
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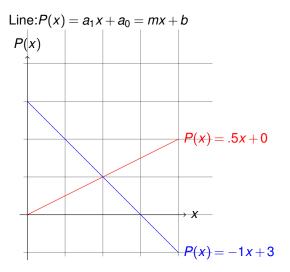
for  $x \in \{0, ..., p-1\}$  and  $a_i \in \{0, ..., p-1\}$ 

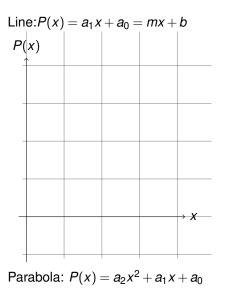
Line:  $P(x) = a_1 x + a_0$ 

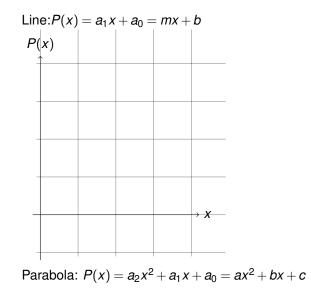
Line: $P(x) = a_1x + a_0 = mx + b$ 

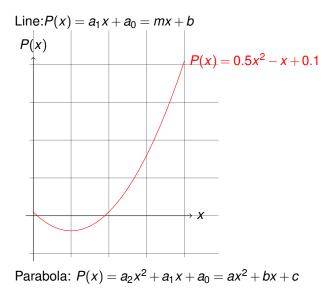


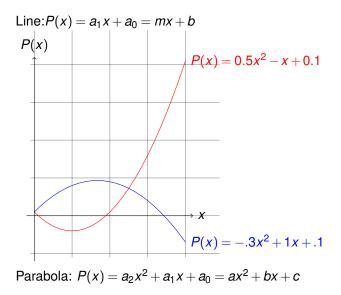




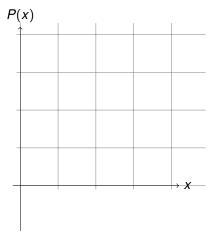




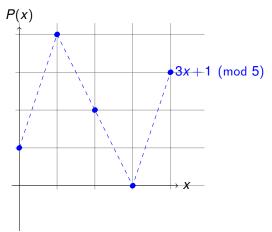




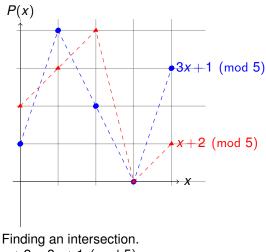
Polynomial:  $P(x) = a_d x^d + \cdots + a_0 \pmod{p}$ 



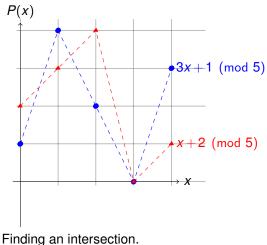
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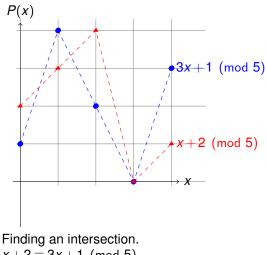
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Finding an intersection.  $x+2 \equiv 3x+1 \pmod{5}$  $\implies 2x \equiv 1 \pmod{5}$  Polynomial:  $P(x) = a_d x^d + \cdots + a_0 \pmod{p}$ 



 $x + 2 \equiv 3x + 1 \pmod{5}$  $\implies 2x \equiv 1 \pmod{5} \implies x \equiv 3 \pmod{5}$ 3 is multiplicative inverse of 2 modulo 5. Polynomial:  $P(x) = a_d x^d + \cdots + a_0 \pmod{p}$ 



 $\begin{array}{l} x+2\equiv 3x+1 \pmod{5} \\ \Longrightarrow 2x\equiv 1 \pmod{5} \implies x\equiv 3 \pmod{5} \\ 3 \text{ is multiplicative inverse of 2 modulo 5.} \\ \text{Good when modulus is prime!!} \end{array}$ 

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**Fact:** Given d + 1 points<sup>1</sup>, exactly 1 degree  $\leq d$  polynomial contains them.

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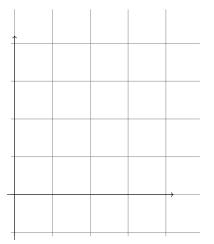
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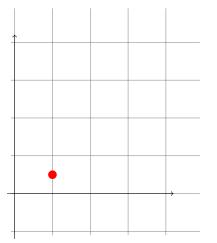
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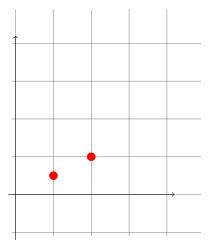
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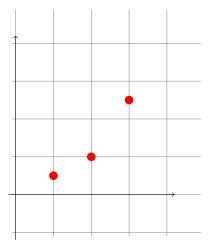
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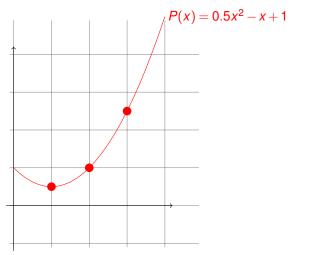
(A) and (D)

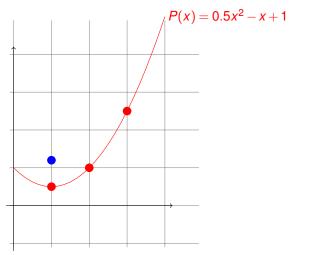


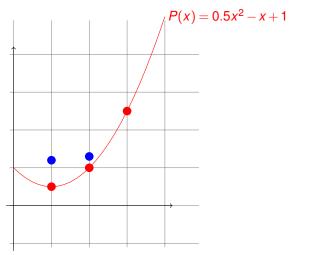


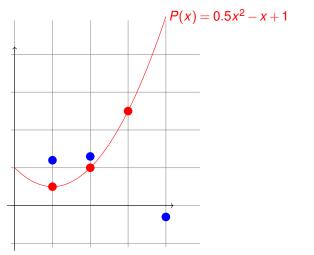


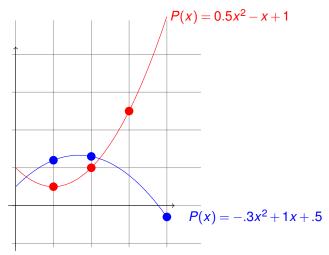




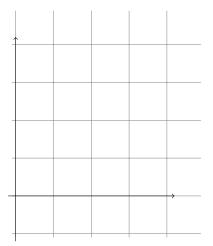


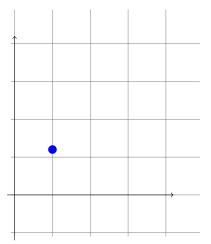


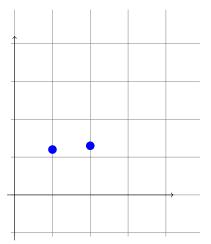


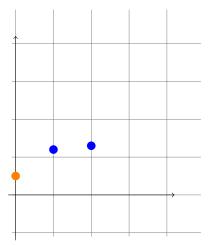


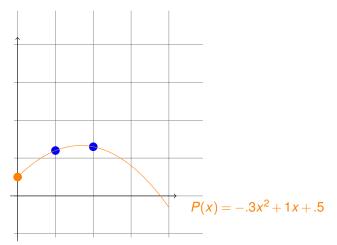
<sup>&</sup>lt;sup>2</sup>Points with different x values.

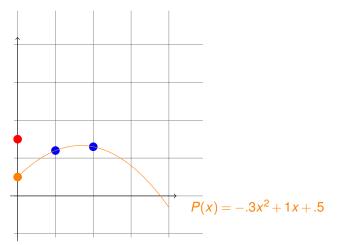


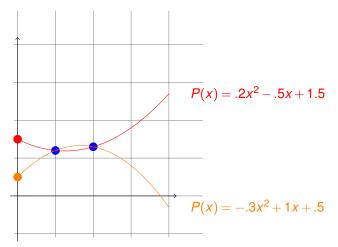


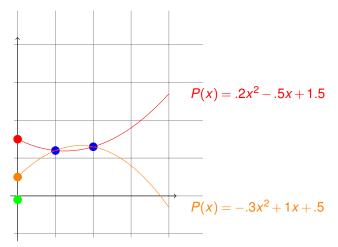


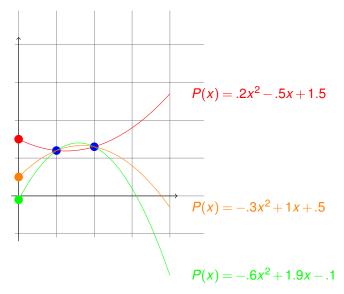


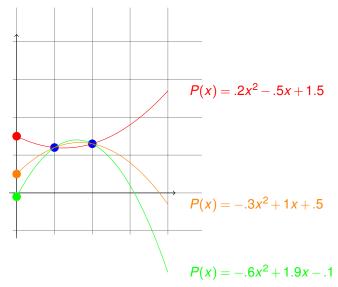












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The polynomial from the scheme:  $P(x) = 2x^2 + 1x + 3 \pmod{5}$ . What is true for the secret sharing scheme using P(x)?

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- (C) A share could be (1,5) because P(1) = 5
- (D) A share could be (2,4)
- (E) A share could be (0,3)

(B), (C) are true. (E) undesirable (reveals secret), start shares from i = 1.

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*P*(1) =

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Backsolve:  $b \equiv 2 \pmod{5}$ . Secret is 2. And the line is...

 $x+2 \mod 5$ .

For a quadratic polynomial,  $a_2x^2 + a_1x + a_0$  hits (1,2); (2,4); (3,0).

$$P(1) = a_2 + a_1 + a_0 \equiv 2 \pmod{5}$$

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So polynomial is  $2x^2 + 1x + 4 \pmod{5}$ 

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**Modular Arithmetic Fact:** Exactly 1 degree  $\leq d$  polynomial with arithmetic modulo prime *p* contains *d* + 1 pts.

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Given d + 1 points, use  $\Delta_i$  functions to go through points?  $(x_1, y_1), \ldots, (x_{d+1}, y_{d+1}).$ Will  $y_1 \Delta_1(x)$  contain  $(x_1, y_1)$ ? Will  $y_2 \Delta_2(x)$  contain  $(x_2, y_2)$ ? Does  $y_1 \Delta_1(x) + y_2 \Delta_2(x)$  contain  $(x_1, y_1)$ ? and  $(x_2, y_2)$ ? See the idea?

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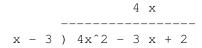
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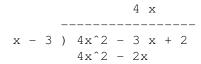
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Must prove Roots fact.





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In general, divide  $P(x)$  by  $(x - a)$  gives  $Q(x)$  and remainder  $r$ .  
That is,  $P(x) = (x - a)Q(x) + r$ 

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**Modular Arithmetic Fact:** Exactly one polynomial degree  $\leq d$  over GF(p), P(x), that hits d+1 points.

#### Shamir's k out of n Scheme:

Secret  $s \in \{0, ..., p-1\}$ 

1. Choose  $a_0 = s$ , and randomly  $a_1, \ldots, a_{k-1}$ .

2. Let 
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**Robustness:** Any *k* knows secret. Knowing *k* pts, only one P(x), evaluate P(0). **Secrecy:** Any k - 1 knows nothing. Knowing  $\leq k - 1$  pts, any P(0) is possible.

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- With k 1 shares, any of p values possible for P(0)!

## Runtime.

#### Runtime.

Runtime: polynomial in k, n, and  $\log p$ .

- 1. Evaluate degree k 1 polynomial *n* times using log *p*-bit numbers.
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Infinite number for reals, rationals, complex numbers!

# Summary

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Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Can solve linear system.

Solution exists: lagrange interpolation.

Unique:

Roots fact: Factoring: (x - r) is root.

Induction only *d* roots.

Apply: P(x), Q(x) degree d.

P(x) - Q(x) is degree  $d \implies d$  roots.

P(x) = Q(x) on d+1 points  $\implies P(x) = Q(x)$ .

Secret Sharing:

k points on degree k - 1 polynomial is great! Can hand out *n* points on polynomial as shares.