Outline for Today.

Polynomials.

Secret Sharing.

Finite Fields.

Share secret among *n* people.

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Secrecy: Any k - 1 knows nothing.

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Two points make a line.

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Two points make a line. Lots of lines go through one point.

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A polynomial

$$P(x)=a_dx^d+a_{d-1}x^{d-1}\cdots+a_0.$$

is specified by **coefficients** $a_d, \ldots a_0$.

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is specified by **coefficients** $a_d, \dots a_0$. P(x) **contains** point (a, b) if b = P(a). **Polynomials over reals**: $a_1, \dots, a_d \in \Re$, use $x \in \Re$.

E.g., Reals, rationals, complex numbers.

E.g., Reals, rationals, complex numbers. Not E.g., the integers.

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Intuitively, a field is a set with operations corresponding to addition, multiplication, and division.



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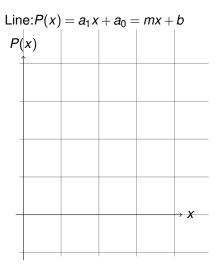
Polynomials P(x) with arithmetic modulo p:

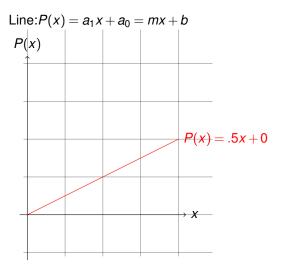
$$P(x) = a_d x^d + a_{d-1} x^{d-1} \cdots + a_0 \pmod{p},$$

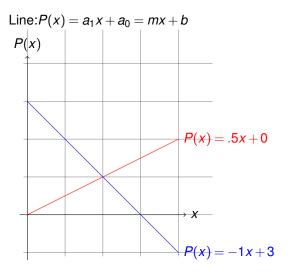
for $x \in \{0, ..., p-1\}$ and $a_i \in \{0, ..., p-1\}$

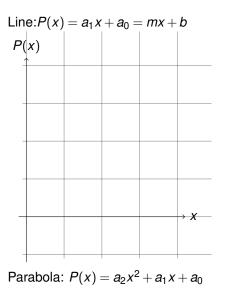
Line: $P(x) = a_1 x + a_0$

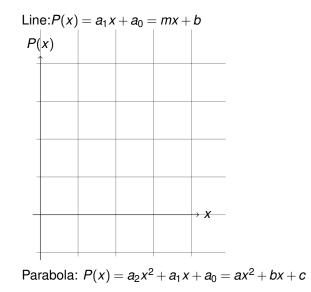
Line: $P(x) = a_1x + a_0 = mx + b$

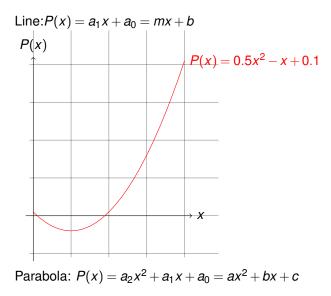


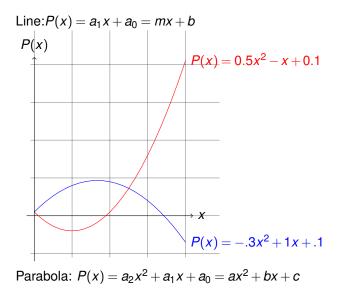




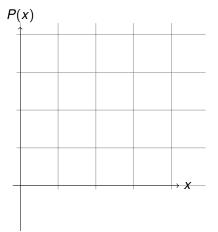




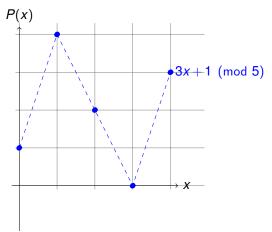




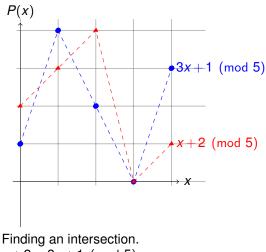
Polynomial: $P(x) = a_d x^d + \cdots + a_0 \pmod{p}$



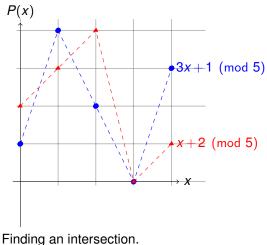
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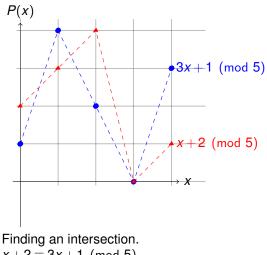
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Finding an intersection. $x+2 \equiv 3x+1 \pmod{5}$ $\implies 2x \equiv 1 \pmod{5}$ Polynomial: $P(x) = a_d x^d + \cdots + a_0 \pmod{p}$



 $x + 2 \equiv 3x + 1 \pmod{5}$ $\implies 2x \equiv 1 \pmod{5} \implies x \equiv 3 \pmod{5}$ 3 is multiplicative inverse of 2 modulo 5. Polynomial: $P(x) = a_d x^d + \cdots + a_0 \pmod{p}$



 $\begin{array}{l} x+2\equiv 3x+1 \pmod{5} \\ \Longrightarrow 2x\equiv 1 \pmod{5} \implies x\equiv 3 \pmod{5} \\ 3 \text{ is multiplicative inverse of 2 modulo 5.} \\ \text{Good when modulus is prime!!} \end{array}$

Two points make a line.

Fact: Given d + 1 points¹, exactly 1 degree $\leq d$ polynomial contains them.

¹Points with different x values.

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Polynomial: $a_n x^n + \cdots + a_0$.

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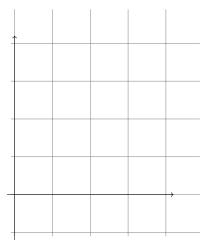
(A) $a_1 = m$ (B) $a_1 = b$ (C) $a_0 = m$ (D) $a_0 = b$.

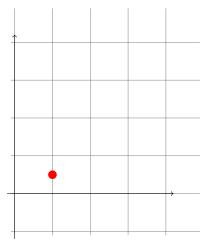
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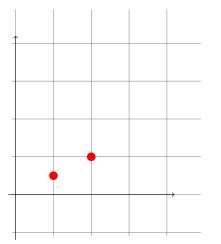
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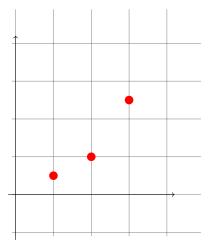
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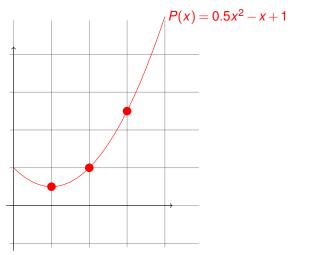
(A) and (D)

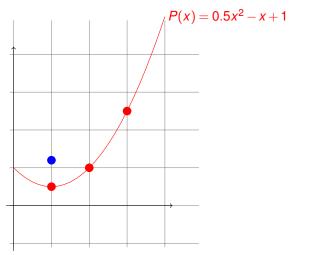


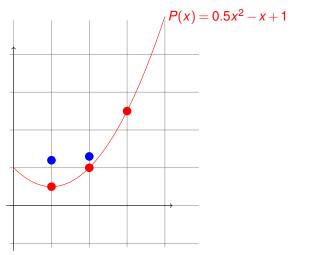


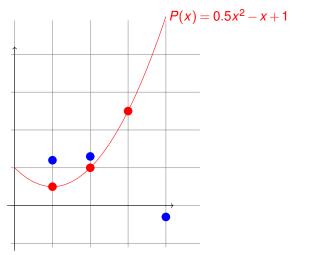


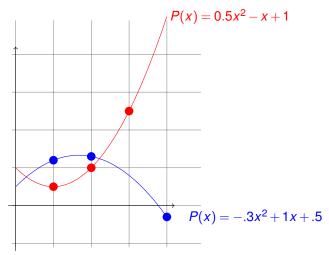




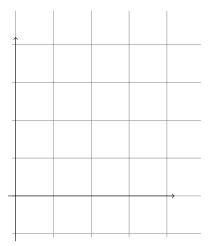


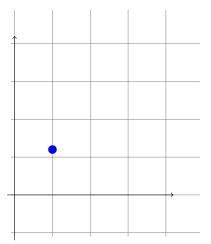


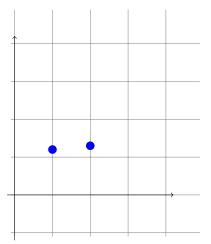


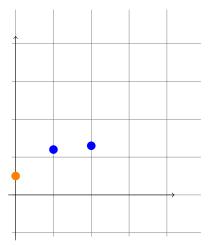


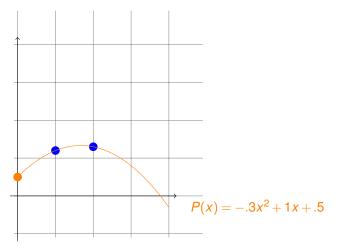
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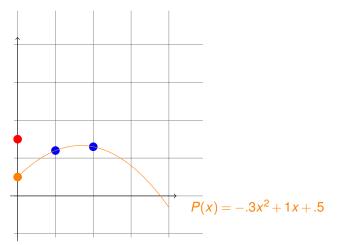


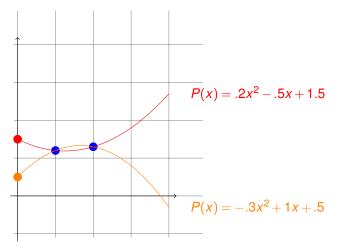


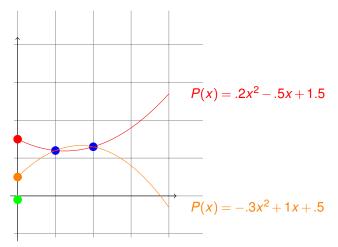


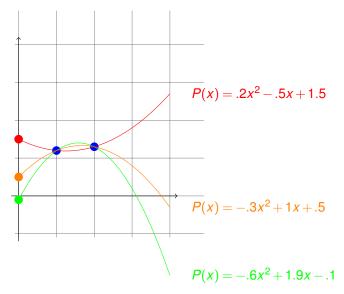


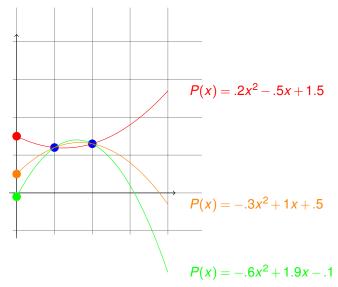












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- (B) The secret is "3".
- (C) A share could be (1,5) because P(1) = 5
- (D) A share could be (2,4)
- (E) A share could be (0,3)

(B), (C) are true. (E) undesirable (reveals secret), start shares from i = 1.

For a line, $a_1x + a_0 = mx + b$ contains points (1,3) and (2,4).

P(1) =

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Backsolve: $b \equiv 2 \pmod{5}$. Secret is 2. And the line is...

 $x+2 \mod 5$.

For a quadratic polynomial, $a_2x^2 + a_1x + a_0$ hits (1,2); (2,4); (3,0).

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$a_2 + a_1 + a_0$	\equiv	2	(mod 5)
$3a_1 + 2a_0$	\equiv	1	(mod 5)
4 <i>a</i> ₁ +2 <i>a</i> ₀	\equiv	2	(mod 5)

For a quadratic polynomial, $a_2x^2 + a_1x + a_0$ hits (1,2); (2,4); (3,0). Plug in points to find equations.

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$a_2 + a_1 + a_0$	\equiv	2	(mod 5)
3 <i>a</i> ₁ +2 <i>a</i> ₀	≡	1	(mod 5)
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So polynomial is $2x^2 + 1x + 4 \pmod{5}$

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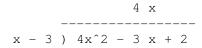
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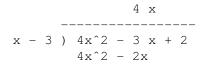
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Must prove Roots fact.





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Divide $4x^2 - 3x + 2$ by (x - 3) modulo 5.

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In general, divide $P(x)$ by $(x - a)$ gives $Q(x)$ and remainder r .
That is, $P(x) = (x - a)Q(x) + r$

Only *d* roots.

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Modular Arithmetic Fact: Exactly one polynomial degree $\leq d$ over GF(p), P(x), that hits d+1 points.

Shamir's k out of n Scheme:

Secret $s \in \{0, ..., p-1\}$

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2. Let
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Robustness: Any *k* knows secret. Knowing *k* pts, only one P(x), evaluate P(0). **Secrecy:** Any k - 1 knows nothing. Knowing $\leq k - 1$ pts, any P(0) is possible.

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- With k shares, reconstruct polynomial, P(x).
- With k 1 shares, any of p values possible for P(0)!

Runtime.

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Runtime: polynomial in k, n, and $\log p$.

- 1. Evaluate degree k 1 polynomial *n* times using log *p*-bit numbers.
- 2. Reconstruct secret by solving system of *k* equations using log *p*-bit arithmetic.

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Infinite number for reals, rationals, complex numbers!

Summary

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Compute solution: *m*, *b*.

Unique:

Assume two solutions, show they are the same.

Today: d + 1 points make a unique degree d polynomial.

Can solve linear system.

Solution exists: lagrange interpolation.

Unique:

Roots fact: Factoring: (x - r) is root.

Induction only *d* roots.

Apply: P(x), Q(x) degree d.

P(x) - Q(x) is degree $d \implies d$ roots.

P(x) = Q(x) on d+1 points $\implies P(x) = Q(x)$.

Secret Sharing:

k points on degree k - 1 polynomial is great! Can hand out *n* points on polynomial as shares.