CS 70 Discrete Mathematics and Probability Theory Summer 2017 Lu, Moulos, and Tang DIS 2B

1 Graph Basics

In the first few parts, you will be answering questions on the following graph G.



- (a) What are the vertex and edge sets V and E for graph G?
- (b) Which vertex has the highest in-degree? Which vertex has the lowest in-degree? Which vertices have the same in-degree and out-degree?
- (c) What are the paths from vertex *B* to *F*, assuming no vertex is visited twice? Which one is the shortest path?
- (d) Which of the following are cycles in G?
 - i. {(B,C), (C,D), (D,B)}
 ii. {(F,G), (G,F)}
 iii. {(A,B), (B,C), (C,D), (D,B)}
 iv. {(B,C), (C,D), (D,H), (H,G), (G,F), (F,E), (E,D), (D,B)}
- (e) Which of the following are walks in *G*?

i.
$$\{(E,G)\}$$

ii. $\{(E,G), (G,F)\}$
iii. $\{(F,G), (G,F)\}$
iv. $\{(A,B), (B,C), (C,D)\}$
v. $\{(E,G), (G,F), (F,G), (G,F)\}$
vi. $\{(E,D), (D,B), (B,E), (E,D), (D,H), (H,G), (G,F)\}$

- (f) Which of the following are tours in G?
 - i. $\{(E,G)\}$ ii. $\{(E,G), (G,F)\}$ iii. $\{(F,G), (G,F)\}$ iv. $\{(E,D), (D,B), (B,E), (E,D), (D,H), (H,G), (G,F)\}$

In the following three parts, let's consider a general undirected graph G with n vertices $(n \ge 3)$.

- (g) True/False: If each vertex of G has degree at most 1, then G does not have a cycle.
- (h) True/False: If each vertex of G has degree at least 2, then G has a cycle.
- (i) True/False: If each vertex of G has degree at most 2, then G is not connected.

2 Build-Up Error?

What is wrong with the following "proof"?

False Claim: If every vertex in an undirected graph has degree at least 1, then the graph is connected.

Proof: We use induction on the number of vertices $n \ge 1$.

Base case: There is only one graph with a single vertex and it has degree 0. Therefore, the base case is vacuously true, since the if-part is false.

Inductive hypothesis: Assume the claim is true for some $n \ge 1$.

Inductive step: We prove the claim is also true for n + 1. Consider an undirected graph on n vertices in which every vertex has degree at least 1. By the inductive hypothesis, this graph is connected. Now add one more vertex x to obtain a graph on (n + 1) vertices, as shown below.



All that remains is to check that there is a path from x to every other vertex z. Since x has degree at least 1, there is an edge from x to some other vertex; call it y. Thus, we can obtain a path from x to z by adjoining the edge $\{x, y\}$ to the path from y to z. This proves the claim for n + 1.

3 Bipartite Graph

Consider an undirected bipartite graph with two disjoint sets L, R. Prove that a graph is bipartite if and only if it no cycles of odd length.