

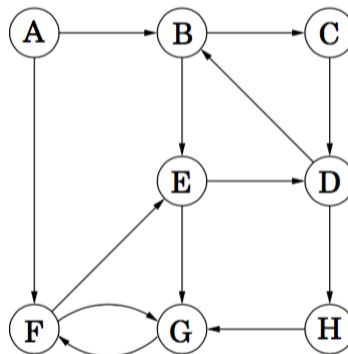
Due: Friday 9/17, 10:00 PM  
Grace period until Friday 9/17 11:59 PM

## Sundry

Before you start writing your final homework submission, state briefly how you worked on it. Who else did you work with? List names and email addresses. (In case of homework party, you can just describe the group.)

## 1 Graph Basics

In the first few parts, you will be answering questions on the following graph  $G$ .



- What are the vertex and edge sets  $V$  and  $E$  for graph  $G$ ?
- Which vertex has the highest in-degree? Which vertex has the lowest in-degree? Which vertices have the same in-degree and out-degree?
- What are the paths from vertex  $B$  to  $F$ , assuming no vertex is visited twice? Which one is the shortest path?
- Which of the following are cycles in  $G$ ?
  - $(B,C), (C,D), (D,B)$
  - $(F,G), (G,F)$
  - $(A,B), (B,C), (C,D), (D,B)$

iv.  $(B,C), (C,D), (D,H), (H,G), (G,F), (F,E), (E,D), (D,B)$

(e) Which of the following are walks in  $G$ ?

i.  $(E, G)$

ii.  $(E, G), (G, F)$

iii.  $(F, G), (G, F)$

iv.  $(A, B), (B, C), (C, D), (H, G)$

v.  $(E, G), (G, F), (F, G), (G, C)$

vi.  $(E, D), (D, B), (B, E), (E, D), (D, H), (H, G), (G, F)$

(f) Which of the following are tours in  $G$ ?

i.  $(E, G)$

ii.  $(E, G), (G, F)$

iii.  $(F, G), (G, F)$

iv.  $(E, D), (D, B), (B, E), (E, D), (D, H), (H, G), (G, F)$

v.  $(B, C), (C, D), (D, H), (H, G), (G, F), (F, E), (E, D), (D, B)$

**In the following three parts, let's consider a general undirected graph  $G$  with  $n$  vertices ( $n \geq 3$ ). If true, provide a short proof. If false, show a counterexample.**

(g) True/False: If each vertex of  $G$  has degree at most 1, then  $G$  does not have a cycle.

(h) True/False: If each vertex of  $G$  has degree at least 2, then  $G$  has a cycle.

(i) True/False: If each vertex of  $G$  has degree at most 2, then  $G$  is not connected.

## 2 Bipartite Graphs

An undirected graph is bipartite if its vertices can be partitioned into two disjoint sets  $L, R$  such that each edge connects a vertex in  $L$  to a vertex in  $R$  (so there does not exist an edge that connects two vertices in  $L$  or two vertices in  $R$ ).

(a) Suppose that a graph  $G$  is bipartite, with  $L$  and  $R$  being a bipartite partition of the vertices. Prove that  $\sum_{v \in L} \deg(v) = \sum_{v \in R} \deg(v)$ .

(b) Suppose that a graph  $G$  is bipartite, with  $L$  and  $R$  being a bipartite partition of the vertices. Let  $s$  and  $t$  denote the average degree of vertices in  $L$  and  $R$  respectively. Prove that  $s/t = |R|/|L|$ .

(c) Prove that a graph is bipartite if and only if it can be 2-colored. (A graph can be 2-colored if every vertex can be assigned one of two colors such that no two adjacent vertices have the same color).

### 3 Touring Hypercube

In the lecture, you have seen that if  $G$  is a hypercube of dimension  $n$ , then

- The vertices of  $G$  are the binary strings of length  $n$ .
- $u$  and  $v$  are connected by an edge if they differ in exactly one bit location.

A *Hamiltonian tour* of a graph is a sequence of vertices  $v_0, v_1, \dots, v_k$  such that:

- Each vertex appears exactly once in the sequence.
- Each pair of consecutive vertices is connected by an edge.
- $v_0$  and  $v_k$  are connected by an edge.

- (a) Show that a hypercube has an Eulerian tour if and only if  $n$  is even. (*Hint: Euler's theorem*)
- (b) Show that every hypercube has a Hamiltonian tour.

### 4 Tournament

A *tournament* is defined to be a directed graph such that for every pair of distinct nodes  $v$  and  $w$ , exactly one of  $(v, w)$  and  $(w, v)$  is an edge (representing which player beat the other in a round-robin tournament). Prove that every tournament has a Hamiltonian path. In other words, you can always arrange the players in a line so that each player beats the next player in the line.

### 5 Planarity and Graph Complements

Let  $G = (V, E)$  be an undirected graph. We define the complement of  $G$  as  $\bar{G} = (V, \bar{E})$  where  $\bar{E} = \{(i, j) | i, j \in V, i \neq j\} - E$ ; that is,  $\bar{G}$  has the same set of vertices as  $G$ , but an edge  $e$  exists in  $\bar{G}$  if and only if it does not exist in  $G$ .

- (a) Suppose  $G$  has  $v$  vertices and  $e$  edges. How many edges does  $\bar{G}$  have?
- (b) Prove that for any graph with at least 13 vertices,  $G$  being planar implies that  $\bar{G}$  is non-planar.
- (c) Now consider the converse of the previous part, i.e., for any graph  $G$  with at least 13 vertices, if  $\bar{G}$  is non-planar, then  $G$  is planar. Construct a counterexample to show that the converse does not hold.

*Hint: Recall that if a graph contains a copy of  $K_5$ , then it is non-planar. Can this fact be used to construct a counterexample?*

## 6 Build-Up Error?

What is wrong with the following "proof"? In addition to finding a counterexample, you should explain what is fundamentally wrong with this approach, and why it demonstrates the danger build-up error.

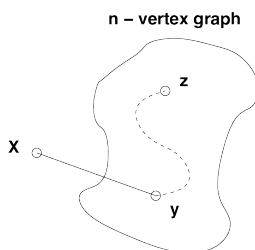
**False Claim:** If every vertex in an undirected graph has degree at least 1, then the graph is connected.

*Proof:* We use induction on the number of vertices  $n \geq 1$ .

*Base case:* There is only one graph with a single vertex and it has degree 0. Therefore, the base case is vacuously true, since the if-part is false.

*Inductive hypothesis:* Assume the claim is true for some  $n \geq 1$ .

*Inductive step:* We prove the claim is also true for  $n + 1$ . Consider an undirected graph on  $n$  vertices in which every vertex has degree at least 1. By the inductive hypothesis, this graph is connected. Now add one more vertex  $x$  to obtain a graph on  $(n + 1)$  vertices, as shown below.



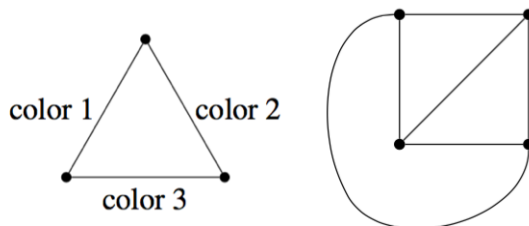
All that remains is to check that there is a path from  $x$  to every other vertex  $z$ . Since  $x$  has degree at least 1, there is an edge from  $x$  to some other vertex; call it  $y$ . Thus, we can obtain a path from  $x$  to  $z$  by adjoining the edge  $\{x, y\}$  to the path from  $y$  to  $z$ . This proves the claim for  $n + 1$ .

## 7 Graph Coloring

Prove that a graph with maximum degree at most  $k$  is  $(k + 1)$ -colorable. (Hint: consider inducting over the number of vertices.)

## 8 Edge Colorings

An edge coloring of a graph is an assignment of colors to edges in a graph where any two edges incident to the same vertex have different colors. An example is shown on the left.



- (a) Show that the 4 vertex complete graph above can be 3 edge colored. (Use the numbers 1,2,3 for colors. A figure is shown on the right.)
- (b) Prove that any graph with maximum degree  $d \geq 1$  can be edge colored with  $2d - 1$  colors.
- (c) Show that a tree can be edge colored with  $d$  colors where  $d$  is the maximum degree of any vertex.