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“Period” divides $p-1$.

Today.

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Secret Sharing.

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Correcting for loss or even corruption.

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Share secret among n people.

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Two points make a line.

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Two points make a line.

Lots of lines go through one point.

Polynomials

A **polynomial**

$$P(x) = a_d x^d + a_{d-1} x^{d-1} \dots + a_0.$$

is specified by **coefficients** $a_d, \dots a_0$.

¹A field is a set of elements with addition and multiplication operations, with inverses. $GF(p) = (\{0, \dots, p-1\}, + \pmod{p}, * \pmod{p})$.

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Note: Often polynomial of degree d means polynomial of at most d .

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Recall polynomial: $a_d x^d + a_{d-1} x^{d-1} + \cdots a_0$.

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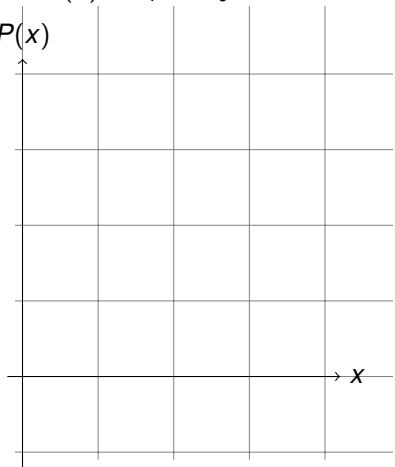
Line: $P(x) = a_1 x + a_0$

Polynomial: $P(x) = a_d x^d + \cdots + a_0$

Line: $P(x) = a_1 x + a_0 = mx + b$

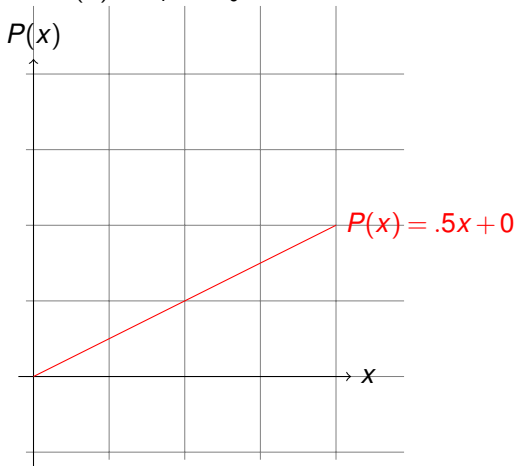
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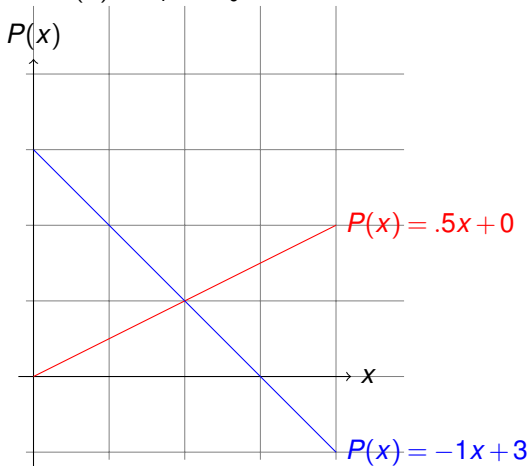
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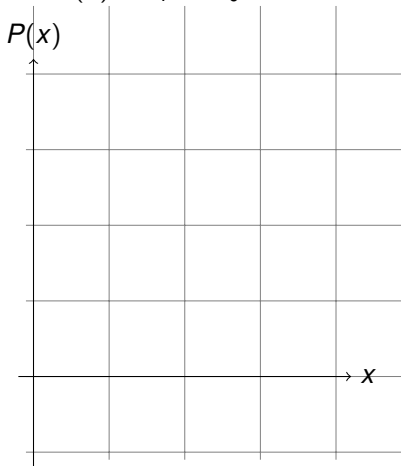
Polynomial: $P(x) = a_d x^4 + \dots + a_0$

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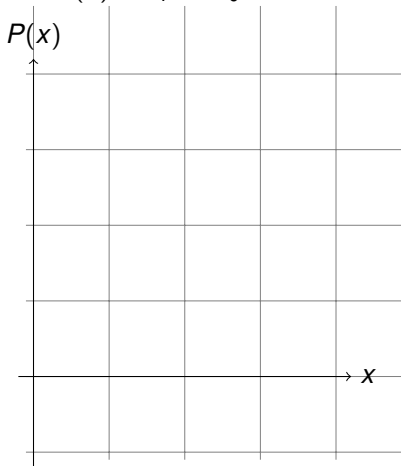
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Parabola: $P(x) = a_2 x^2 + a_1 x + a_0$

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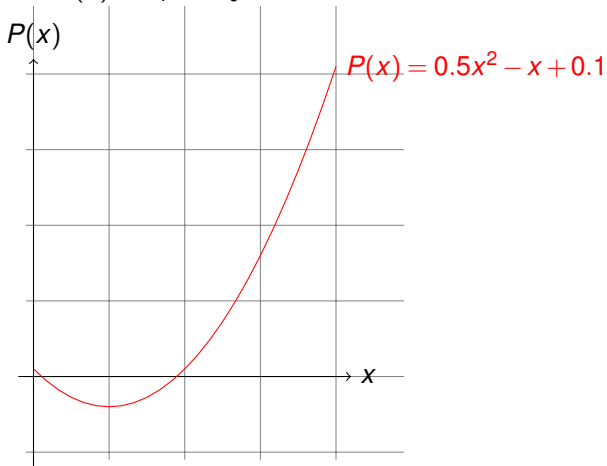
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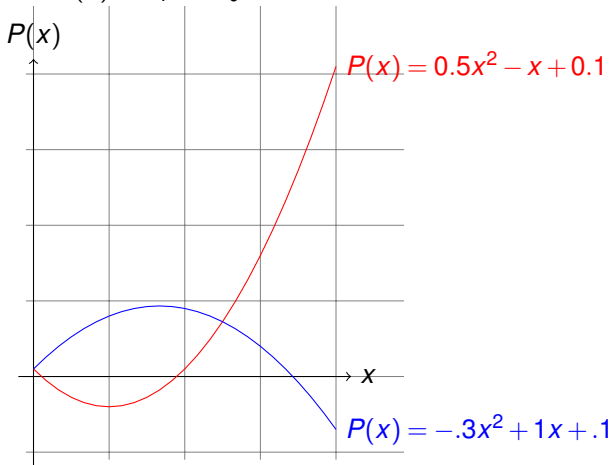
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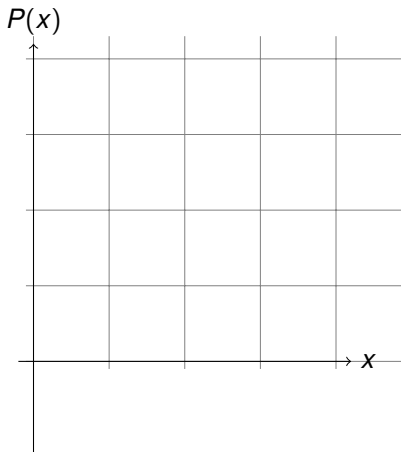
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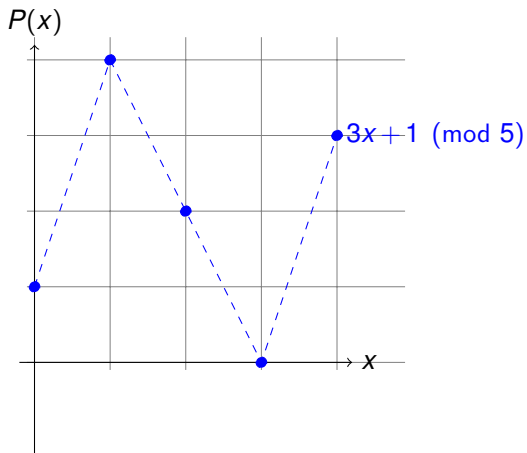


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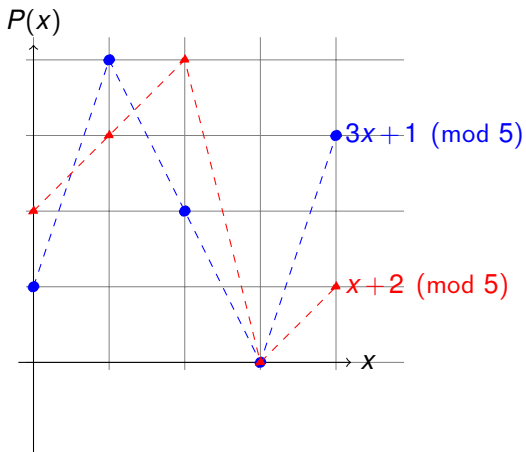
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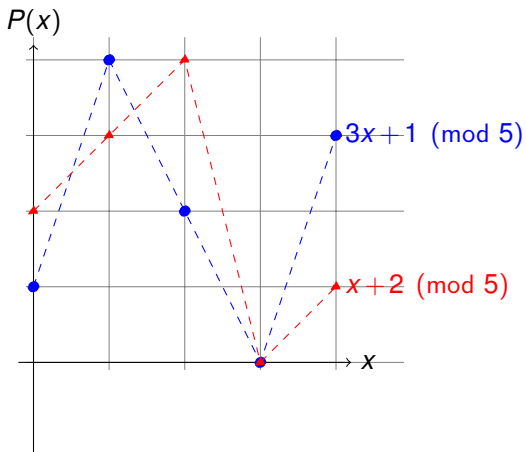


Finding an intersection.

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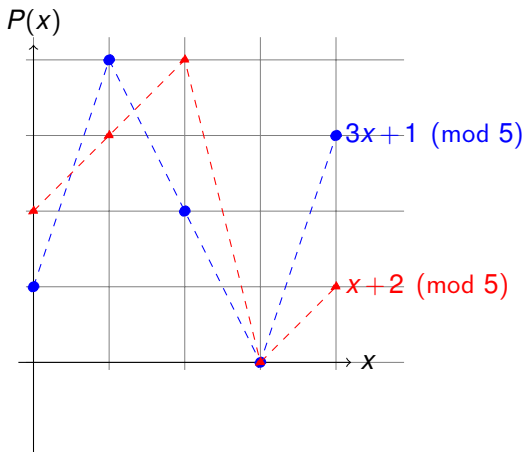
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Good when modulus is prime!!

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Fact: Exactly 1 degree $\leq d$ polynomial contains $d + 1$ points.²

²Points with different x values.

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Poll.

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Say points are $(x_1, y_1), (x_2, y_2)$.

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All true.

In the Flow (Steph Curry) Poll.

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Flow poll. (All true. (B) is not a proof, it is restatement.)

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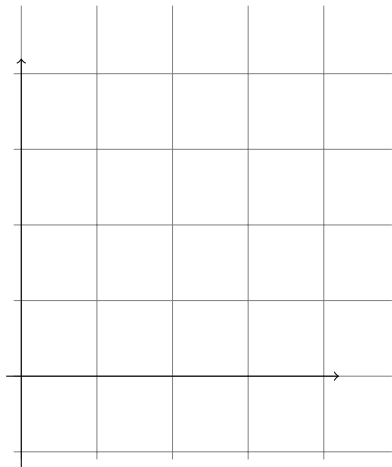
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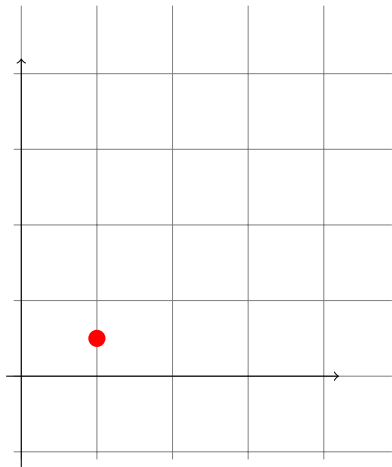
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3 points determine a parabola.



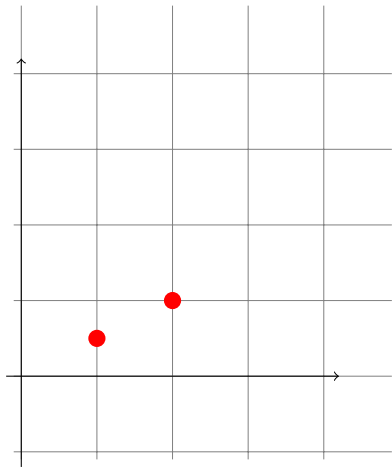
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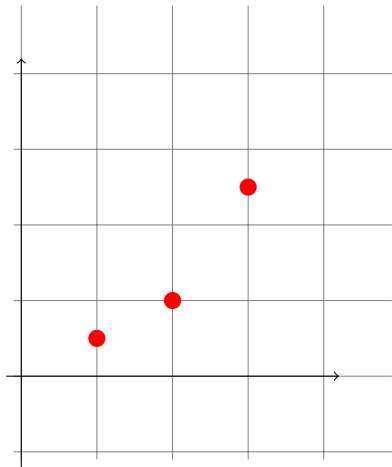
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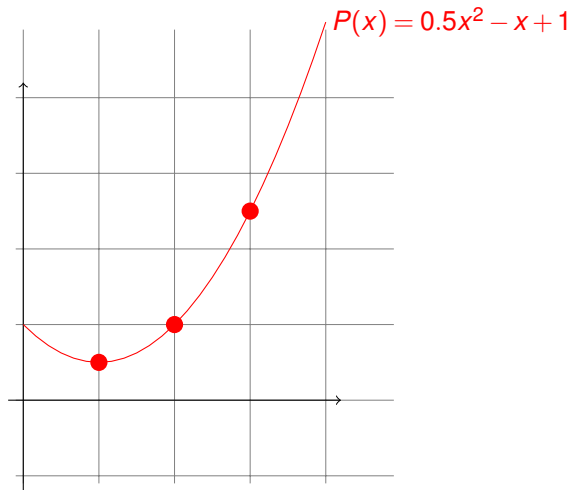
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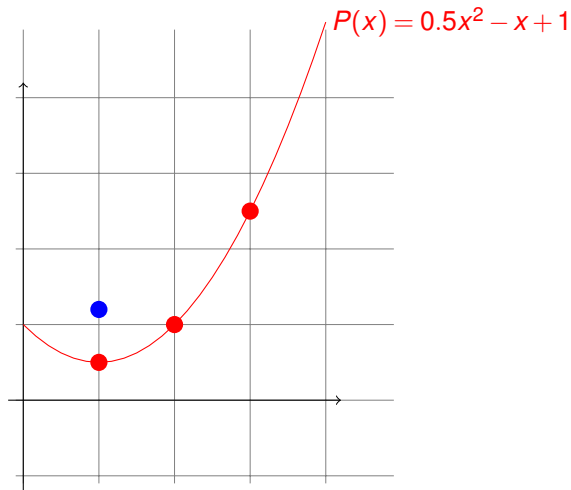
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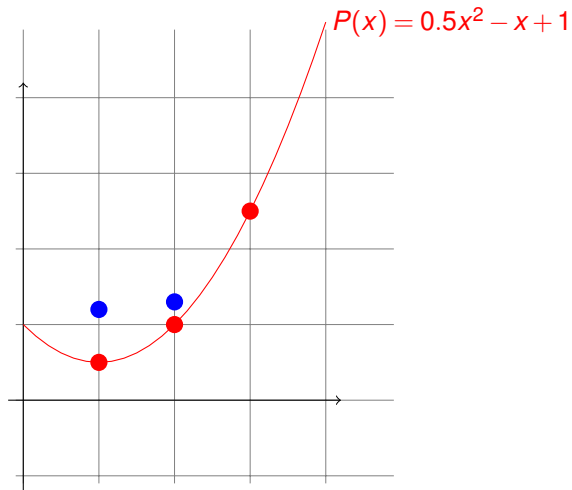
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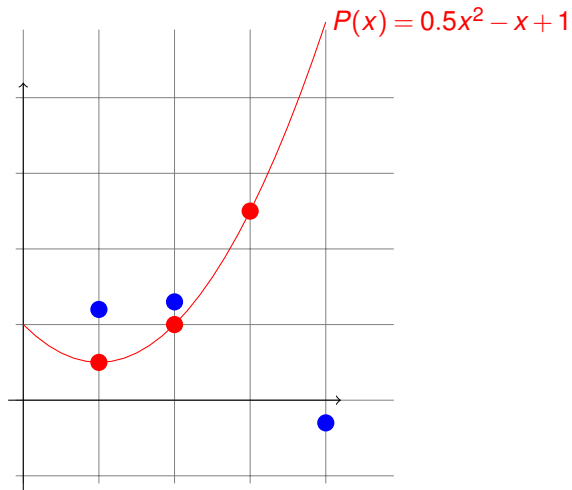
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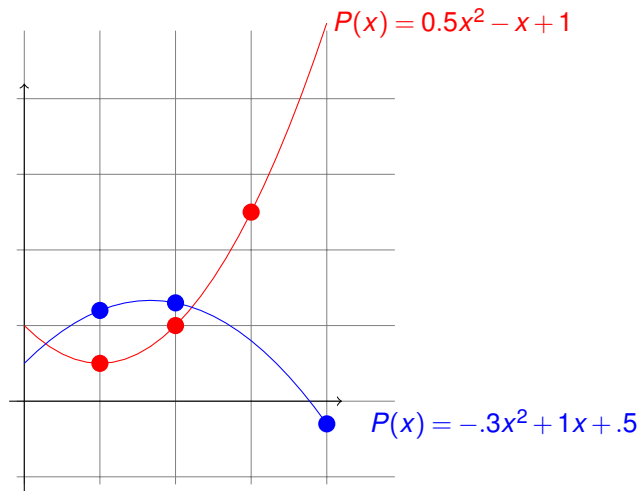
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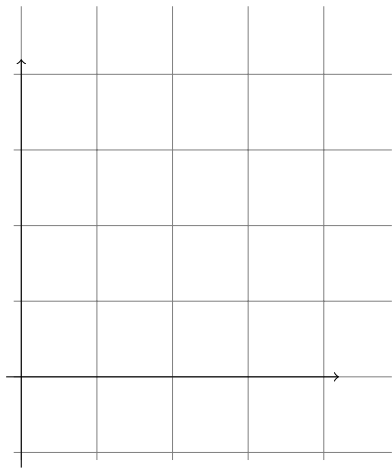
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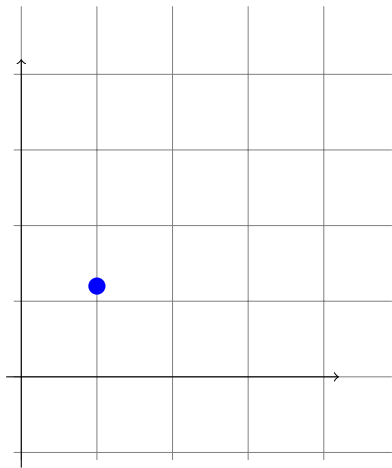
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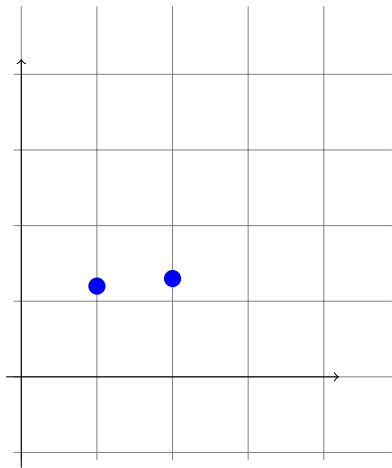
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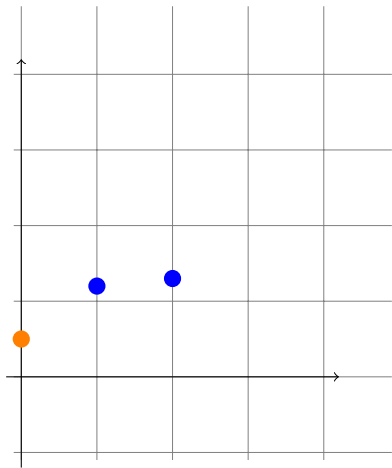
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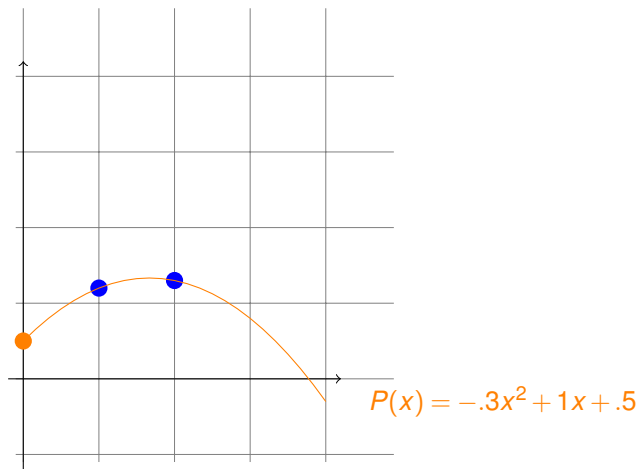
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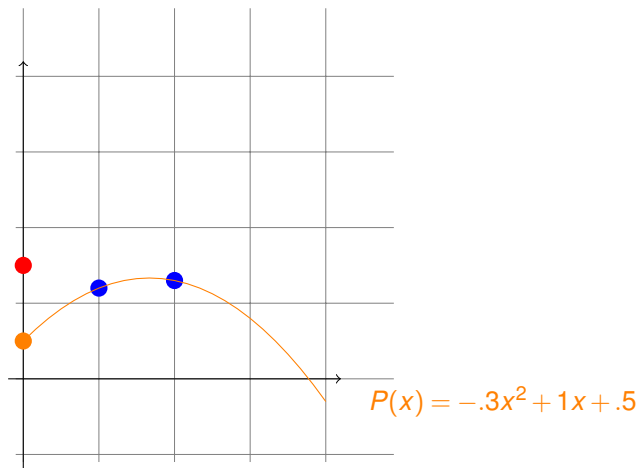
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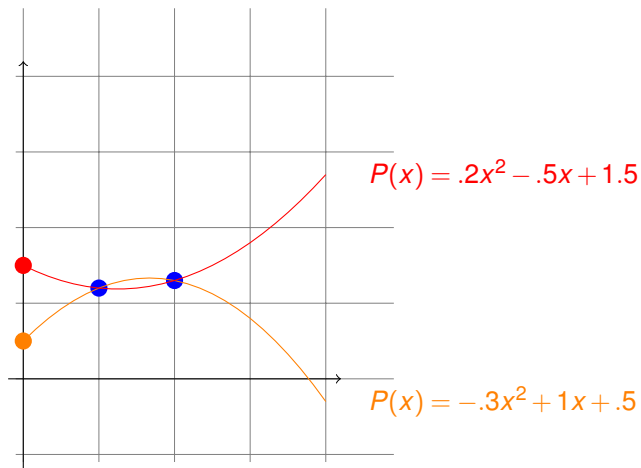
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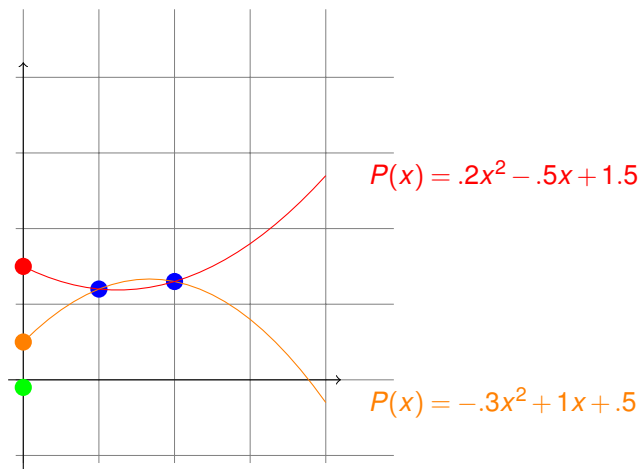
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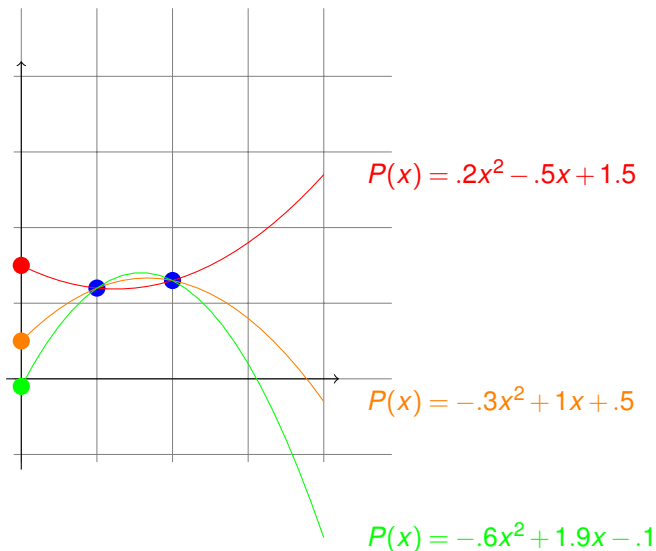
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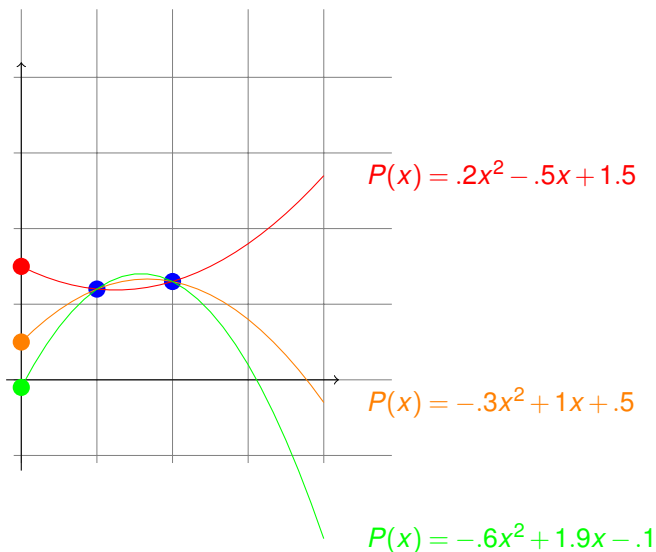


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So polynomial is $2x^2 + 1x + 4 \pmod{5}$

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$$\Delta_i(x_j) = 0 \text{ if } i \neq j \text{ and } \Delta_i(x_i) = 1$$

And..

$$P(x) = y_1 \Delta_1(x) + y_2 \Delta_2(x) + \cdots + y_{d+1} \Delta_{d+1}(x).$$

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And Degree d polynomial.

Construction proves the existence of a polynomial!

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Put the delta functions together.

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Assume two different polynomials $Q(x)$ and $P(x)$ hit the points.

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Must prove **Roots fact**.

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In general, divide $P(x)$ by $(x - a)$ gives $Q(x)$ and remainder r .

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That is, $P(x) = (x - a)Q(x) + r$ where $Q(x)$ has degree $d - 1$.

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Lemma 1: $P(x)$ has root a iff $P(x)/(x - a)$ has remainder 0:

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Plugin a :

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Modular Arithmetic Fact: Exactly one polynomial degree $\leq d$ over $GF(p)$, $P(x)$, that hits $d + 1$ points.

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1. Evaluate degree $k - 1$ polynomial n times using $\log p$ -bit numbers.
2. Reconstruct secret by solving system of k equations using $\log p$ -bit arithmetic.

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$P(x) = Q(x)$ on $d + 1$ points $\implies P(x) = Q(x)$.

Summary

Two points make a line.

Compute solution: m, b .

Unique:

Assume two solutions, show they are the same.

Today: $d + 1$ points make a unique degree d polynomial.

Cuz:

Can solve linear system.

Solution exists: lagrange interpolation.

Unique:

Roots fact: Factoring sez $(x - r)$ is root.

Induction, says only d roots.

Apply: $P(x), Q(x)$ degree d .

$P(x) - Q(x)$ is degree $d \implies d$ roots.

$P(x) = Q(x)$ on $d + 1$ points $\implies P(x) = Q(x)$.

Secret Sharing:

k points on degree $k - 1$ polynomial is great!

Can hand out n points on polynomial as shares.