Today. Variance, covariance. Discuss expectation as predictor. How close to expectation? Using expectation, and variance. What if you predict expectation? Also, prediction from evidence.	Calculating $E[g(X)]$ : LOTUSLet $Y = g(X)$ . Assume that we know the distribution of X.We want to calculate $E[Y]$ .Method 1: We calculate the distribution of Y: $Pr[Y = y] = Pr[X \in g^{-1}(y)]$ where $g^{-1}(x) = \{x \in \Re : g(x) = y\}$ .This is typically rather tedious!Method 2: We use the following result.Called "Law of the unconscious statistician."Theorem: $E[g(X)] = \sum_{x} g(x)Pr[X = x]$ .Proof: $E[g(X)] = \sum_{x} g(x)Pr[\omega] = \sum_{x} \sum_{\omega \in X^{-1}(x)} g(X(\omega))Pr[\omega]$ $= \sum_{x} \sum_{\omega \in X^{-1}(x)} g(x)Pr[\omega] = \sum_{x} g(x)Pr[\omega] = \sum_{x} g(x)Pr[\omega]$ $= \sum_{x} g(x)Pr[X = x]$ .	Poll.Which is LOTUS?(A) $E[X] = \sum_{x \in \text{Range}(X)} g(x) \times Pr[g(X) = g(x)]$ No. overcounts $Pr[g(X) = g(x)]$ .(B) $E[X] = \sum_{x \in \text{Range}(X)} g(x) \times Pr[X = x]$ Yes. May count $g(x)$ twice, if $g(x) = g(x')$ .(C) $E[X] = \sum_{x \in \text{Range}(g)} x \times Pr[g(X) = x]$ No. $g(x)$ is image, x is pre-image.
Geometric Distribution. Experiment: flip a coin with heads prob. <i>p</i> . until Heads. Random Variable <i>X</i> : number of flips. And distribution is: (A) $X \sim G(p) : Pr[X = i] = (1 - p)^{i-1}p.$ (B) $X \sim B(p, n) : Pr[X = i] = {n \choose i}p^i(1 - p)^{n-i}.$ (A) Distribution of $X \sim G(p) : Pr[X = i] = (1 - p)^{i-1}p.$	Geometric Distribution: Memoryless - Interpretation $Pr[X > n+m X > n] = Pr[X > m], m, n \ge 0.$ $\begin{matrix} B & A \\ T \top T \top \dots T \top T \top T \top T \cdots T \\ n & m \end{matrix}$ $Pr[X > n+m X > n] = Pr[A B] = Pr[A'] = Pr[X > m].$ $A': \text{ is } m \text{ coin tosses before heads.}$ $A B: m \text{ 'more' coin tosses before heads.}$ The coin is memoryless, therefore, so is X. Independent coin: $Pr[H \text{'any previous set of coin tosses'}] = p$	Geometric Distribution: Memoryless by derivation. Let X be G(p). Then, for $n \ge 0$ , $Pr[X > n] = Pr[$ first n flips are $T] = (1 - p)^n$ . Theorem $Pr[X > n + m X > n] = Pr[X > m], m, n \ge 0$ . Proof: $Pr[X > n + m X > n] = \frac{Pr[X > n + m \text{ and } X > n]}{Pr[X > n]}$ $= \frac{Pr[X > n + m]}{Pr[X > n]}$ $= \frac{Pr[X > n + m]}{Pr[X > n]}$ $= \frac{(1 - p)^{n+m}}{(1 - p)^n} = (1 - p)^m$ = Pr[X > m].







## Fixed points.

Number of fixed points in a random permutation of *n* items. "Number of student that get homework back."

 $X = X_1 + X_2 \cdots + X_n$ where  $X_i$  is indicator variable for *i*th student getting hw back.

$$E(X^{2}) = \sum_{i} E(X_{i}^{2}) + \sum_{i \neq j} E(X_{i}X_{j}).$$
  
=  $n \times \frac{1}{n} + (n)(n-1) \times \frac{1}{n(n-1)}$   
=  $1 + 1 = 2.$ 

$$\begin{split} E(X_i^2) &= 1 \times Pr[X_i = 1] + 0 \times Pr[X_i = 0] \\ &= \frac{1}{n} \\ E(X_i X_j) &= 1 \times Pr[X_i = 1 \cap X_j = 1] + 0 \times Pr[\text{``anything else''}] \\ &= 1 \times \frac{(n-2)!}{n!} = \frac{1}{n(n-1)} \\ Var(X) &= E(X^2) - (E(X))^2 = 2 - 1 = 1. \end{split}$$

## Properties of variance.

- 1.  $Var(cX) = c^2 Var(X)$ , where c is a constant. Scales by  $c^2$ .
- 2. Var(X+c) = Var(X), where c is a constant. Shifts center.

#### Proof:

$$Var(cX) = E((cX)^{2}) - (E(cX))^{2}$$
  
=  $c^{2}E(X^{2}) - c^{2}(E(X))^{2} = c^{2}(E(X^{2}) - E(X)^{2})^{2}$   
=  $c^{2}Var(X)$   
$$Var(X+c) = E((X+c-E(X+c))^{2})$$
  
=  $E((X+c-E(X)-c)^{2})$   
=  $E((X-E(X))^{2}) = Var(X)$ 

### Poll: fixed points.

What's true? (A)  $X_i$  and  $X_j$  are independent. No. If student *i* gets student *j*'s homework. (B)  $E[X_iX_j] = Pr[X_iX_j = 1]$ Yes. Indicator random variable. (C)  $Pr[X_iX_j] = \frac{(n-2)!}{n!}$ Yes. (n-2)! outcomes where  $X_iX_j = 1$ . (D)  $X_i^2 = X_i$ . Yes.  $1^2 = 1$  and  $0^2 = 1$ ,  $X_i \in \{0, 1\}$ .

## Independent random variables.

Independent: P[X = a, Y = b] = Pr[X = a]Pr[Y = b]Fact: E[XY] = E[X]E[Y] for independent random variables.

$$E[XY] = \sum_{a} \sum_{b} a \times b \times Pr[X = a, Y = b]$$
  
= 
$$\sum_{a} \sum_{b} a \times b \times Pr[X = a]Pr[Y = b]$$
  
= 
$$(\sum_{a} aPr[X = a])(\sum_{b} bPr[Y = b])$$
  
= 
$$E[X]E[Y]$$

Variance: binomial.

$$E[X^{2}] = \sum_{i=0}^{n} i^{2} \times {\binom{n}{i}} p^{i} (1-p)^{n-i}.$$
  
= Really???!!##...

Too hard! Ok.. fine. Let's do something else. Maybe not much easier...but there is a payoff.

Variance of sum of two independent random variables Theorem:

If X and Y are independent, then

Var(X+Y) = Var(X) + Var(Y).

Proof:

Since shifting the random variables does not change their variance, let us subtract their means. That is, we assume that E(X) = 0 and E(Y) = 0.

Then, by independence,

$$E(XY) = E(X)E(Y) = 0.$$

Hence,

 $var(X + Y) = E((X + Y)^2) = E(X^2 + 2XY + Y^2)$ =  $E(X^2) + 2E(XY) + E(Y^2) = E(X^2) + E(Y^2)$ = var(X) + var(Y).



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Flip coin with heads probability p.
X- how many heads?
X_i = \begin{cases} 1 & \text{if } ith flip is heads} \\ 0 & \text{otherwise} \end{cases}
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Variance of Binomial Distribution.

 $E(X_i^2) = 1^2 \times p + 0^2 \times (1-p) = p.$   $Var(X_i) = p - (E(X))^2 = p - p^2 = p(1-p).$   $p = 0 \implies Var(X_i) = 0$   $p = 1 \implies Var(X_i) = 0$   $X = X_1 + X_2 + \dots X_n.$   $X_i \text{ and } X_j \text{ are independent: } Pr[X_i = 1|X_j = 1] = Pr[X_i = 1].$  $Var(X) = Var(X_1 + \dots X_n) = np(1-p).$ 

#### Correlation

**Definition** The correlation of *X*, *Y*, *Cor*(*X*, *Y*) is  $corr(X, Y) : \frac{cov(X, Y)}{\sigma(X)\sigma(Y)}.$ 

Theorem:  $-1 \le corr(X, Y) \le 1$ . **Proof:** Idea:  $(a - b)^2 > 0 \rightarrow a^2 + b^2 \ge 2ab$ . Simple case: E[X] = E[Y] = 0 and  $E[X^2] = E[Y^2] = 1$ . Cov(X, Y) = E[XY].  $E[(X - Y)^2] = E[X^2] + E[Y^2] - 2E[XY] = 2(1 - E[XY]) \ge 0$   $\rightarrow E[XY] \le 1$ .  $E[(X + Y)^2] = E[X^2] + E[Y^2] + 2E[XY] = 2(1 + E[XY]) \ge 0$   $\rightarrow E[XY] \ge -1$ . Shifting and scaling doesn't change correlation.

Poisson Distribution: Variance.

**Definition** Poisson Distribution with parameter  $\lambda > 0$ 

$$X = P(\lambda) \Leftrightarrow \Pr[X = m] = \frac{\lambda^m}{m!} e^{-\lambda}, m \ge 0$$

Mean, Variance? Ugh. Recall that Poission is the limit of the Binomial with  $p = \lambda/n$  as  $n \to \infty$ . Mean:  $pn = \lambda$ Variance:  $p(1-p)n = \lambda - \lambda^2/n \to \lambda$ .  $E(X^2)$ ?  $Var(X) = E(X^2) - (E(X))^2$  or  $E(X^2) = Var(X) + E(X)^2$ .  $E(X^2) = \lambda + \lambda^2$ .

## **Examples of Covariance**



Note that E[X] = 0 and E[Y] = 0 in these examples. Then cov(X, Y) = E[XY].

When cov(X, Y) > 0, the RVs X and Y tend to be large or small together. X and Y are said to be positively correlated.

When cov(X, Y) < 0, when X is larger, Y tends to be smaller. X and Y are said to be negatively correlated.

When cov(X, Y) = 0, we say that X and Y are uncorrelated.



#### Markov's inequality

The inequality is named for Andrey Markov, though in work by Pafnuty Chebyshev. (Sometimes) called Chebyshev's first inequality.

**Theorem** Markov's Inequality Assume  $f : \Re \to [0, \infty)$  is nondecreasing. Then,

$$Pr[X \ge a] \le \frac{E[f(X)]}{f(a)}$$
, for all *a* such that  $f(a) > 0$ .

Proof: Claim:

$$1\{X \ge a\} \le \frac{f(X)}{f(a)}.$$

If X < a, the inequality reads  $0 \le f(x)/f(a)$ , since  $f(\cdot) \ge 0$ . If  $X \ge a$ , it reads  $1 \le f(x)/f(a)$ , since  $f(\cdot)$  is nondecreasing.

Taking the expectation yields the inequality, expectation of an indicator is the probability. and expectation is monotone, e.g., weighted sum of points.

That is,  $\sum_{v} \Pr[X = v] \mathbb{1}\{v \ge a\} \le \sum_{v} \Pr[X = v] \frac{f(v)}{f(a)}$ .

ntuition: 
$$E[f(X)] > f(a)Pr[X > a] = f(a)Pr[X > f(a)]$$

## Properties of Covariance cov(X, Y) = E[(X - E[X])(Y - E[Y])] = E[XY] - E[X]E[Y].Fact (a) var[X] = cov(X, X)(b) X, Y independent $\Rightarrow cov(X, Y) = 0$ (c) cov(a + X, b + Y) = cov(X, Y)(d) $cov(aX + bY, cU + dV) = ac \cdot cov(X, U) + ad \cdot cov(X, V)$ $+bc \cdot cov(Y, U) + bd \cdot cov(Y, V).$ Proof: (a)-(b)-(c) are obvious. (d) In view of (c), one can subtract the means and assume that the RVs are zero-mean. Then,

cov(aX + bY, cU + dV) = E[(aX + bY)(cU + dV)]=  $ac \cdot E[XU] + ad \cdot E[XV] + bc \cdot E[YU] + bd \cdot E[YV]$ =  $ac \cdot cov(X, U) + ad \cdot cov(X, V) + bc \cdot cov(Y, U) + bd \cdot cov(Y, V).$ 

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## Lake Woebegone: Poll

#### What is true?

(A) Everyone is above average (on midterm) False. Average would be higher.
(B) For a random variable, at most half the people can be more than twice the average. False. Consder Pr[X = -2] = 1/3 and Pr[X = 1] = 2/3. E[X] = 0.

(C) For the midterm with no negative scores, at most half the people can be more than twice the average.

True. Otherwise average would be higher.

# Markov Inequality Example: G(p)

Let 
$$X = G(p)$$
. Recall that  $E[X] = \frac{1}{p}$  and  $E[X^2] = \frac{2-p}{p^2}$ .





# This is Pafnuty's inequality: Theorem: $Pr[|X-E[X]| > a] \leq \frac{var[X]}{r^2}$ , for all a > 0. **Proof:** Let Y = |X - E[X]| and $f(y) = y^2$ . Then, $\Pr[Y \ge a] \le \frac{E[f(Y)]}{f(a)} = \frac{var[X]}{a^2}.$ This result confirms that the variance measures the "deviations from the mean." Estimation: Expectation and Mean Squared Error. "Best" guess about Y, is E[Y]. If "best" is Mean Squared Error. More precisely, the value of *a* that minimizes $E[(Y - a)^2]$ is a = E[Y]. Proof: Let $\hat{Y} := Y - E[Y]$ . Then, $E[\hat{Y}] = E[Y - E[Y]] = E[Y] - E[Y] = 0.$ So, $E[\hat{Y}c] = 0, \forall c$ . Now, $E[(Y-a)^2] = E[(Y-E[Y]+E[Y]-a)^2]$ $= E[(\hat{Y}+c)^2]$ with c = E[Y] - a $= E[\hat{Y}^{2} + 2\hat{Y}c + c^{2}] = E[\hat{Y}^{2}] + 2E[\hat{Y}c] + c^{2}$ $= E[\hat{Y}^2] + 0 + c^2 \ge E[\hat{Y}^2].$ Hence, $E[(Y - a)^2] \ge E[(Y - E[Y])^2], \forall a$ .

Chebyshev's Inequality



