

CS70 – SPRING 2026

LECTURE 5 : FEB. 3

Today : Graphs

Goals

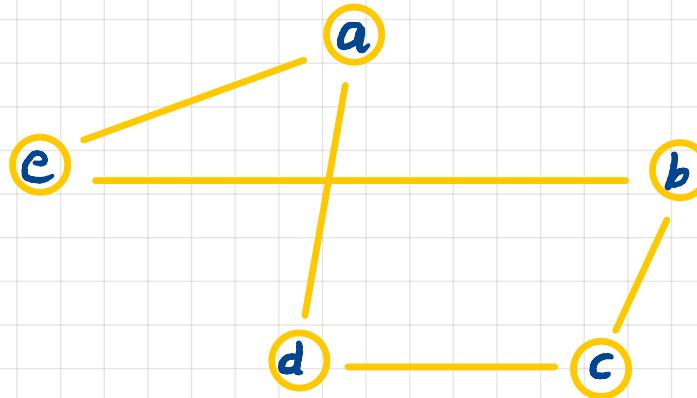
- Understand graphs as an important model for many phenomena in CS & real world
- Explore some basic properties of graphs :
 - paths in graphs
 - trees & complete graphs
 - planar graphs
 - connectivity & hypercubes

} next
lecture

Graphs

$G = (V, E)$

vertices edges

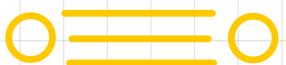


$$V = \{a, b, c, d, e\}$$

$$E = \{ \{a, e\}, \{a, d\}, \{e, b\}, \{b, c\}, \{c, d\} \}$$

An edge is an unordered pair of vertices

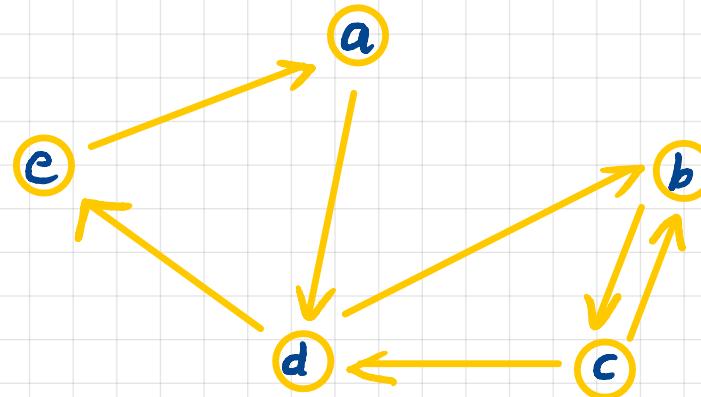
Unless otherwise stated, we don't allow loops ?
or multiple edges



i.e., our graphs are simple graphs, not multigraphs

Directed Graphs

$$G = (V, E)$$



$$V = \{a, b, c, d, e\}$$

$$E = \{(e, a), (d, e), (a, d), (d, b), (b, c), (c, b), (c, d)\}$$

Same, except edges are now ordered pairs of vertices

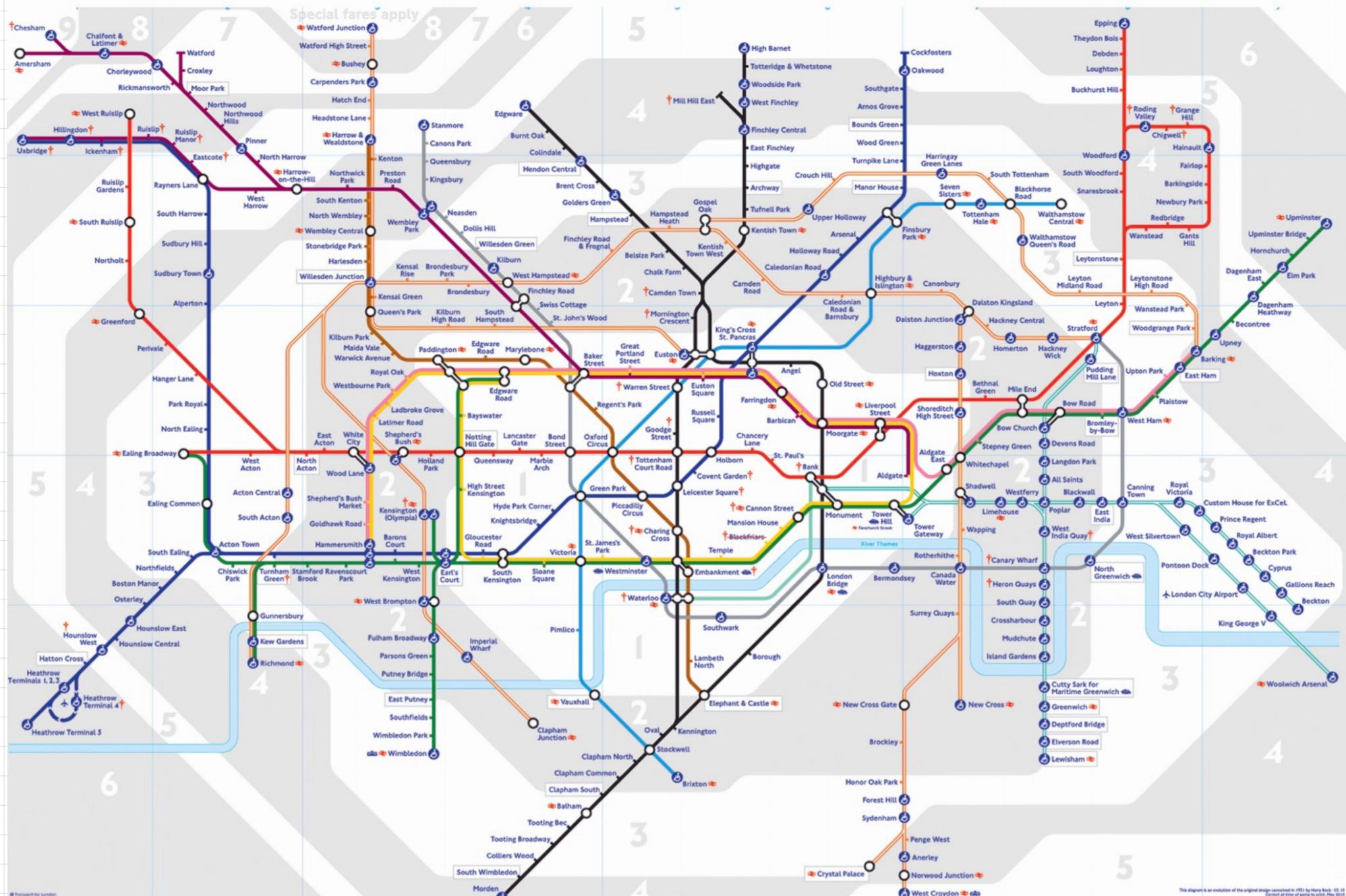
Note : Undirected graphs are a special case of directed graphs :

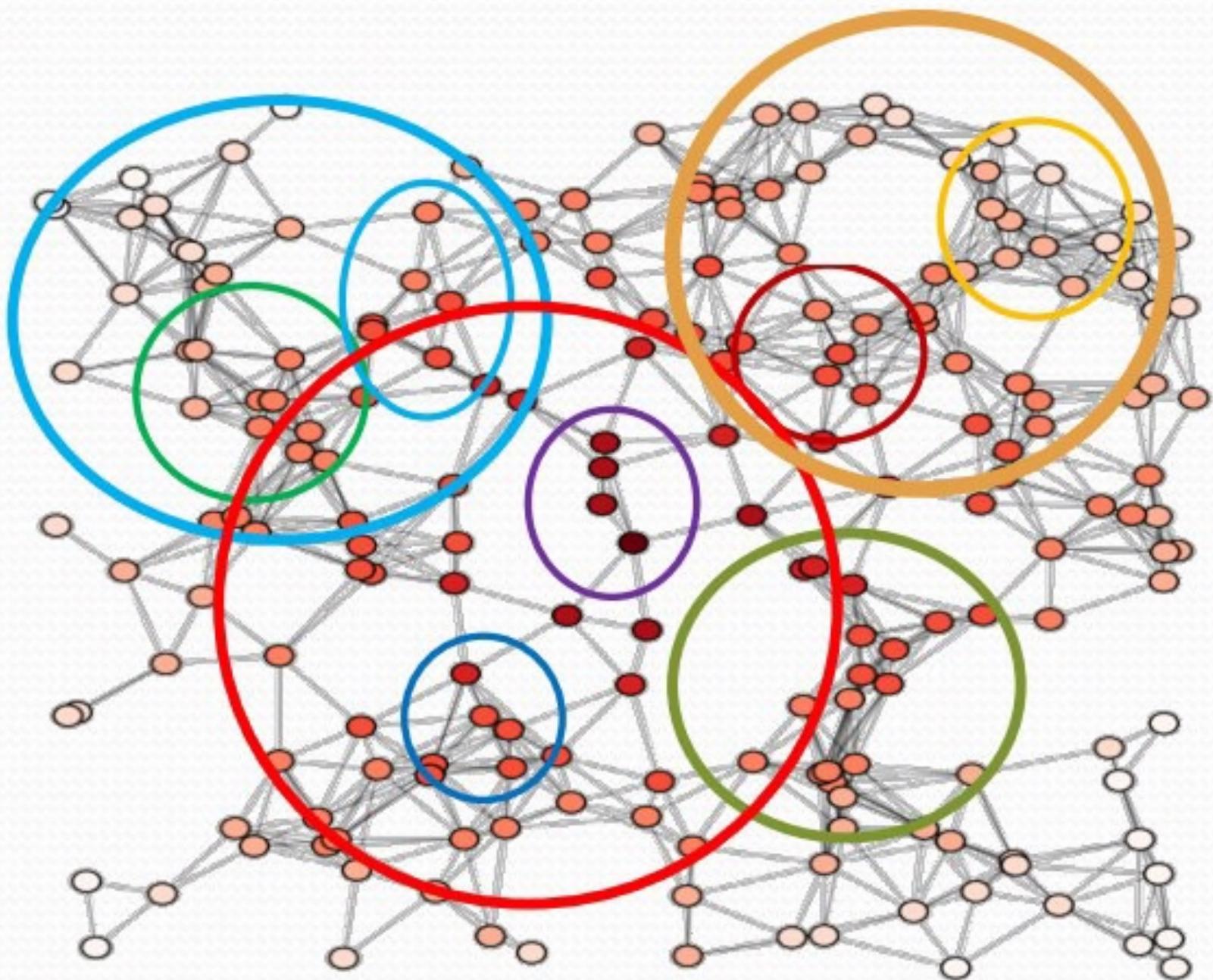
$$\text{---} \equiv \overleftarrow{\text{---}} \overrightarrow{\text{---}}$$

Examples of Undirected Graphs

- road networks (no one-way streets)
- power grids
- social networks ("knows")

: : :



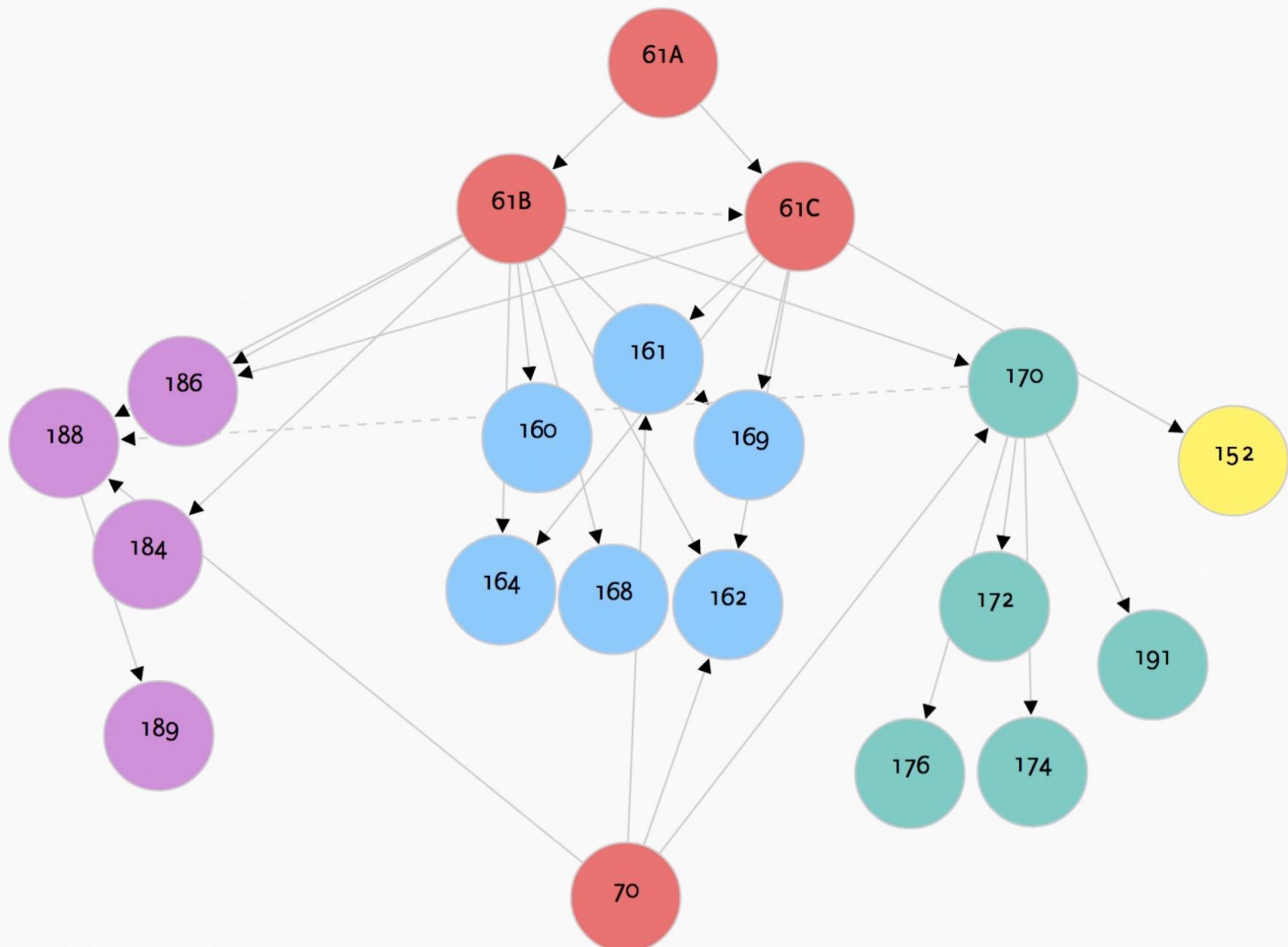


Renren Social Network: Chen Wen Xin

Examples of Directed Graphs

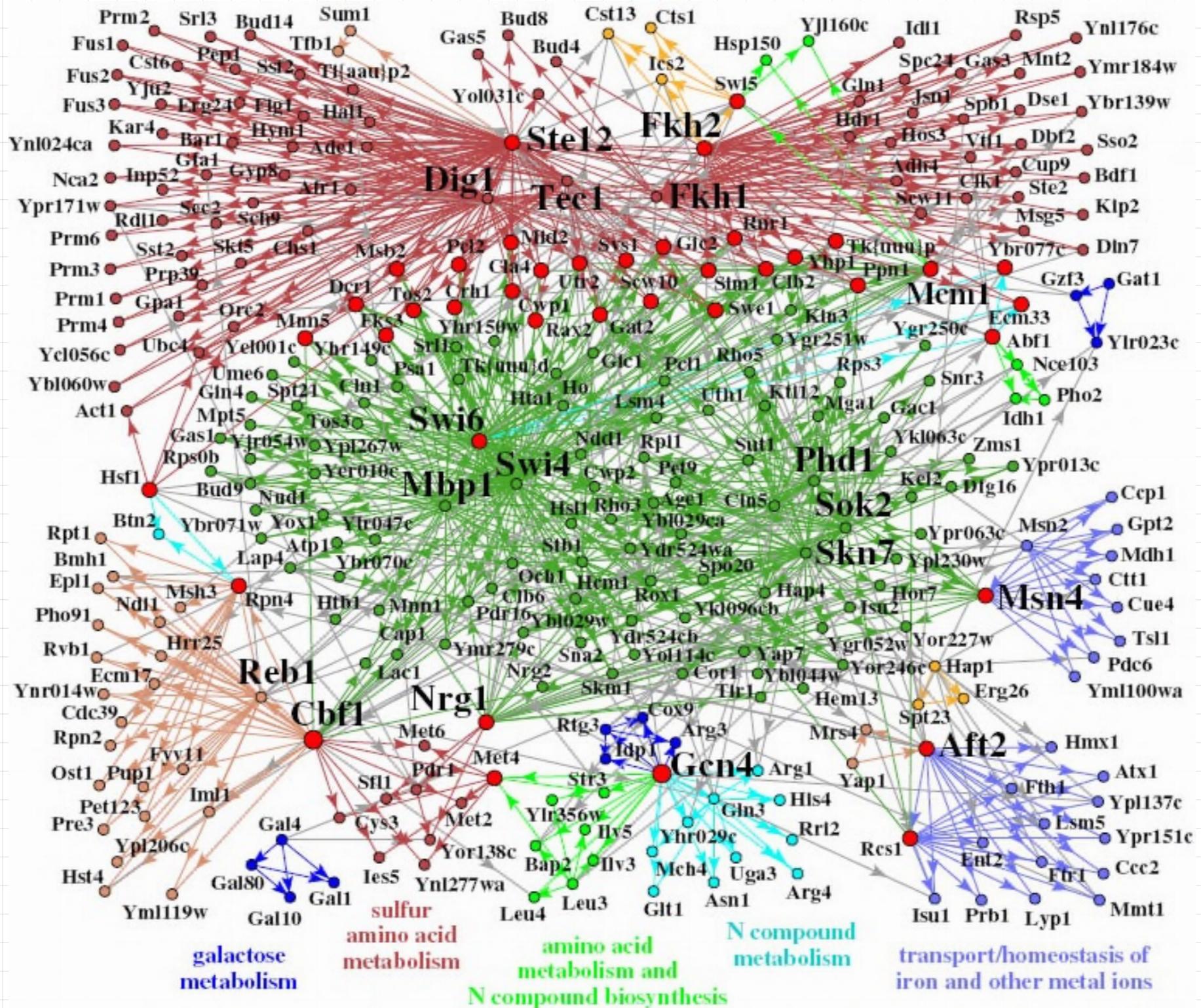
- internet (web links)
- course prerequisites
- biological networks
- social networks ("recognizes")

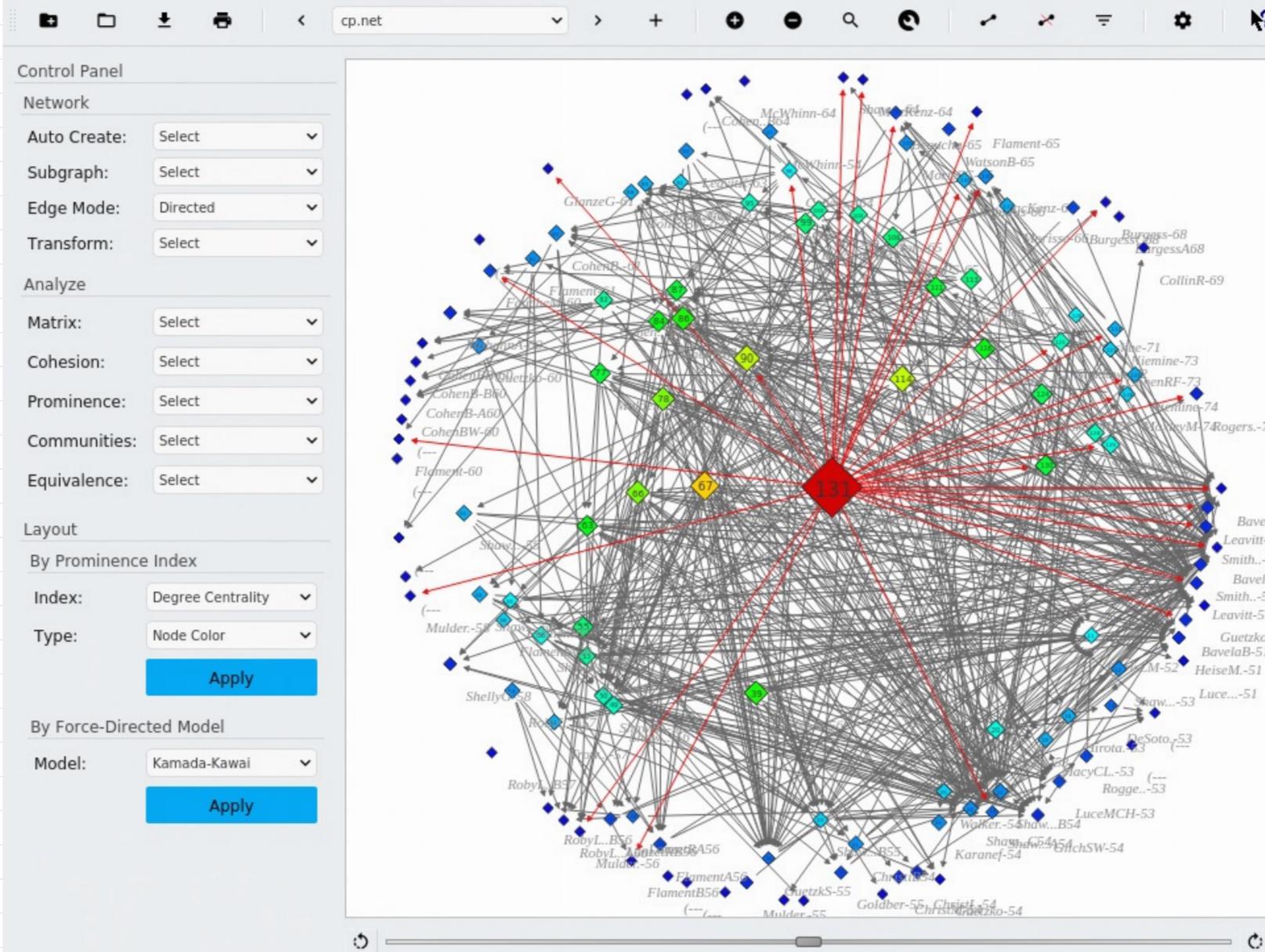
⋮



Required →

Recommended ➤





Statistics Panel

Network

Type: Directed
 Nodes: 131
 Arcs: 631
 Density: 0.0370523

Selection

Nodes: 1
 Arcs: 0

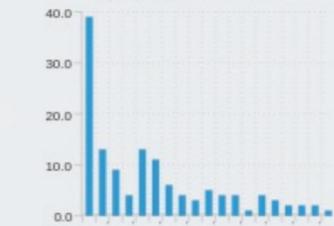
Clicked Node

Number: 131
 In-Degree: 0
 Out-Degree: 29
 Clu.Coeff.: 0

Clicked Edge

Name: -
 Weight: -

(out)Degree distribution



Some Terminology



$e = \{u, v\}$ edge

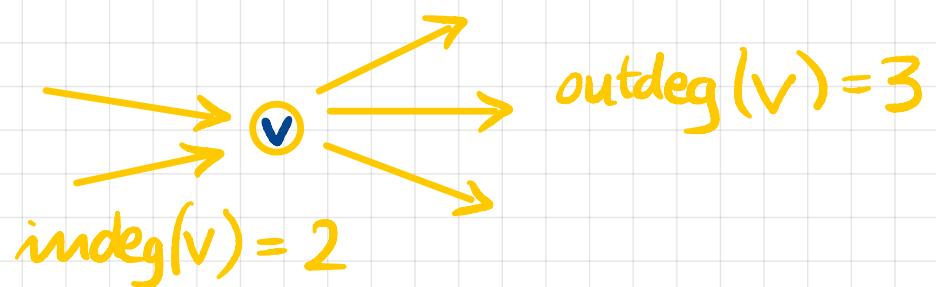
u, v are adjacent or neighbors
e is incident on u & v



degree of v = number of neighbors

If v has degree 0 it is isolated

For directed graphs :



Q: What is the sum of all vertex degrees in G ?

A: $\sum_{v \in V} \deg(v) =$

More Terminology

A (simple) path is a sequence of edges

$$\{v_1, v_2\}, \{v_2, v_3\}, \{v_3, v_4\}, \dots, \{v_{k-1}, v_k\}$$

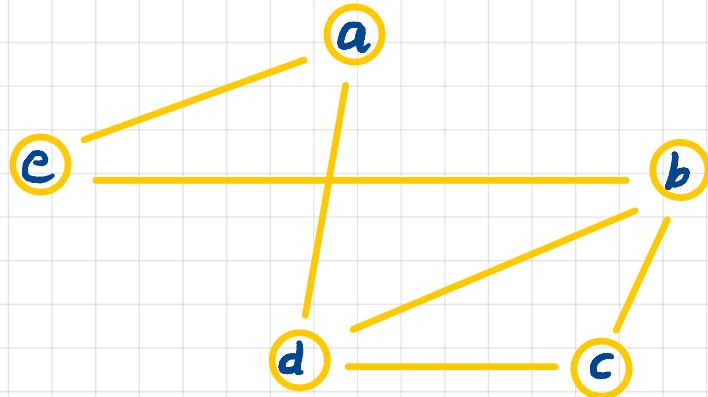
s.t. all the v_i are distinct (except possibly $v_k = v_1$)

[Equivalently: $v_1 - v_2 - v_3 - \dots - v_k$]

A cycle is a path where $v_1 = v_k$

A walk is any sequence of edges as above (repeated vertices/edges allowed)

A tour is a walk where $v_1 = v_k$



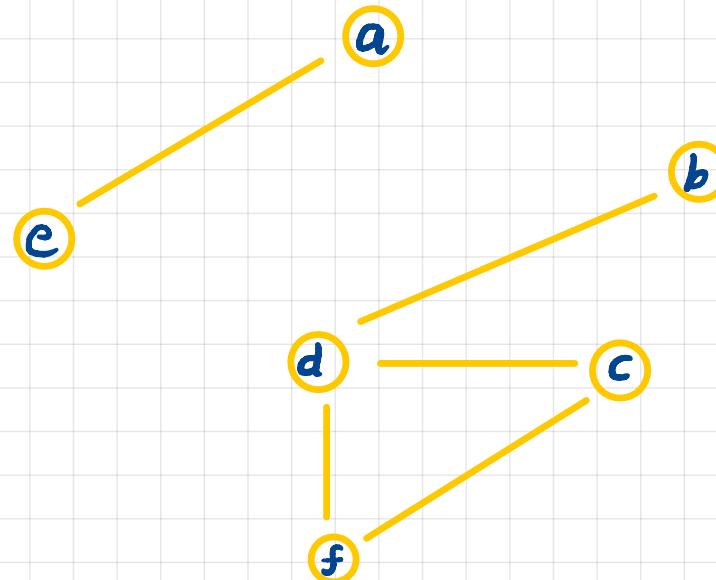
E.g.

path: $e - a - d - c$

cycle: $d - c - b - d$

A graph is connected if there is a path from any vertex to any other vertex

Any graph can be decomposed into connected components

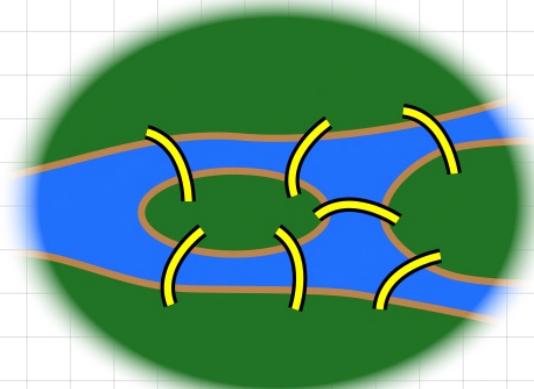
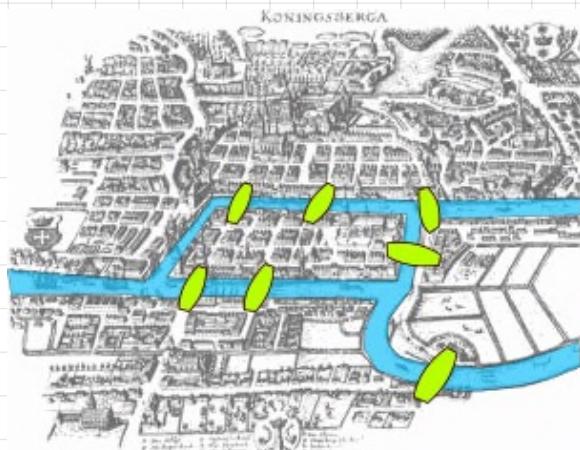


Graph Theory : Topics

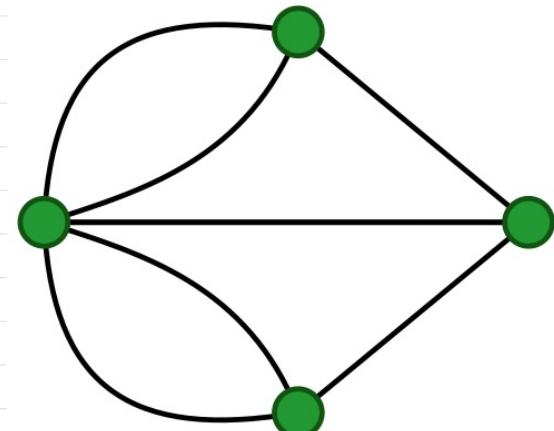
- Eulerian tours
- Trees & complete graphs
- Planar graphs
- Connectivity & hypercubes

} next lecture

Euler : Bridges of Königsberg (1736)



Can we cross all
7 bridges once
& end up where
we started ?



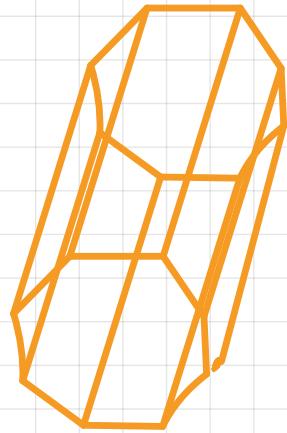
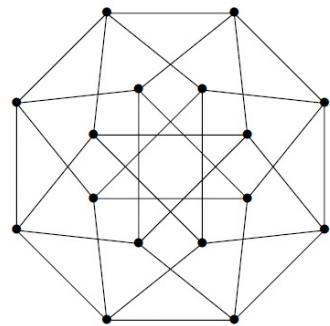
(multi)graph



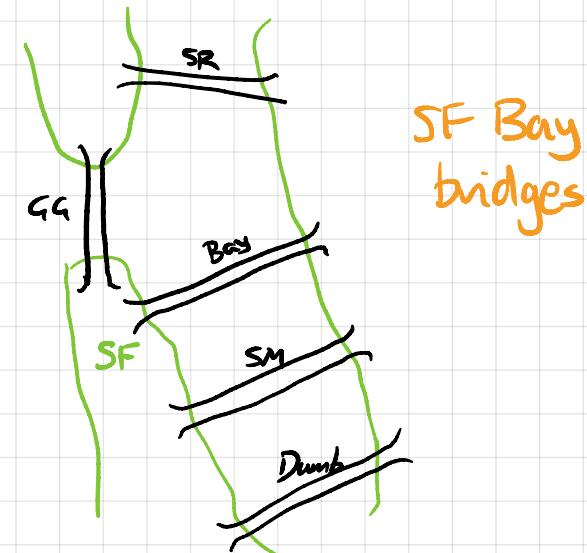
Defn : An Eulerian tour of a graph is a walk that traverses each edge exactly once and ends up at the starting vertex

Theorem [Euler] : A connected graph has an Eulerian tour \Leftrightarrow the degree of every vertex is even

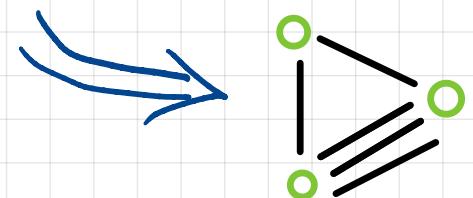
Examples



octagonal
"lampshade"



SF Bay
bridges



Theorem [Euler] : A connected graph has an Eulerian tour \Leftrightarrow the degree of every vertex is even

Proof :

subroutine FindTour(G, s)

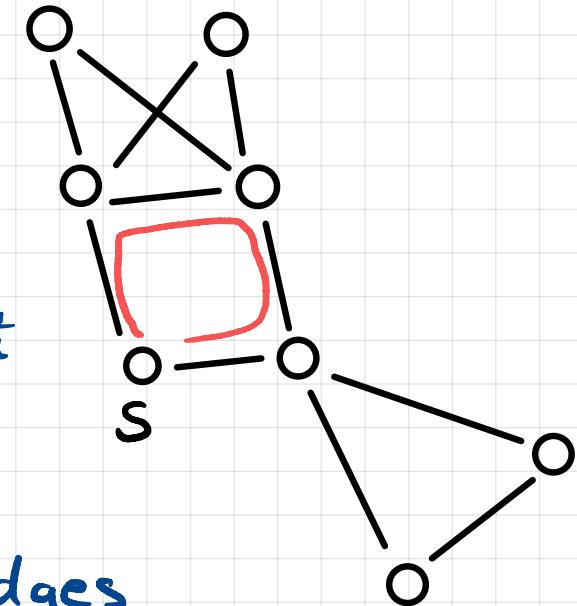
start at s

repeat

choose any untraversed edge incident
on current vertex & traverse it

until get stuck

return the tour formed by traversed edges



Claim: $\text{FindTour}(G, s)$ always gets stuck at s

Proof:

Euler(G, s)

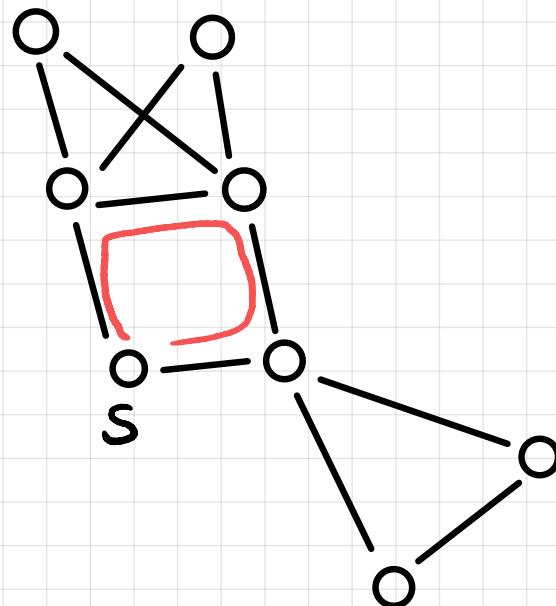
$T = \text{FindTour}(G, s)$

remove edges of T from G

let $\{G_1, \dots, G_k\}$ be the connected components

of the remaining graph, and s_i the first vertex of G_i visited by T

output($\text{Splice}(T, \text{Euler}(G_1, s_1), \dots, \text{Euler}(G_k, s_k))$)



Theorem : For any connected G with even degrees, $\text{Euler}(G, s)$ outputs an Euler tour starting & ending at s

Proof : (Strong) induction on # edges m in G

Base : $m=0$ ✓ (nothing to prove)

Ind. step :

Eulerian Tours – Wrap-Up

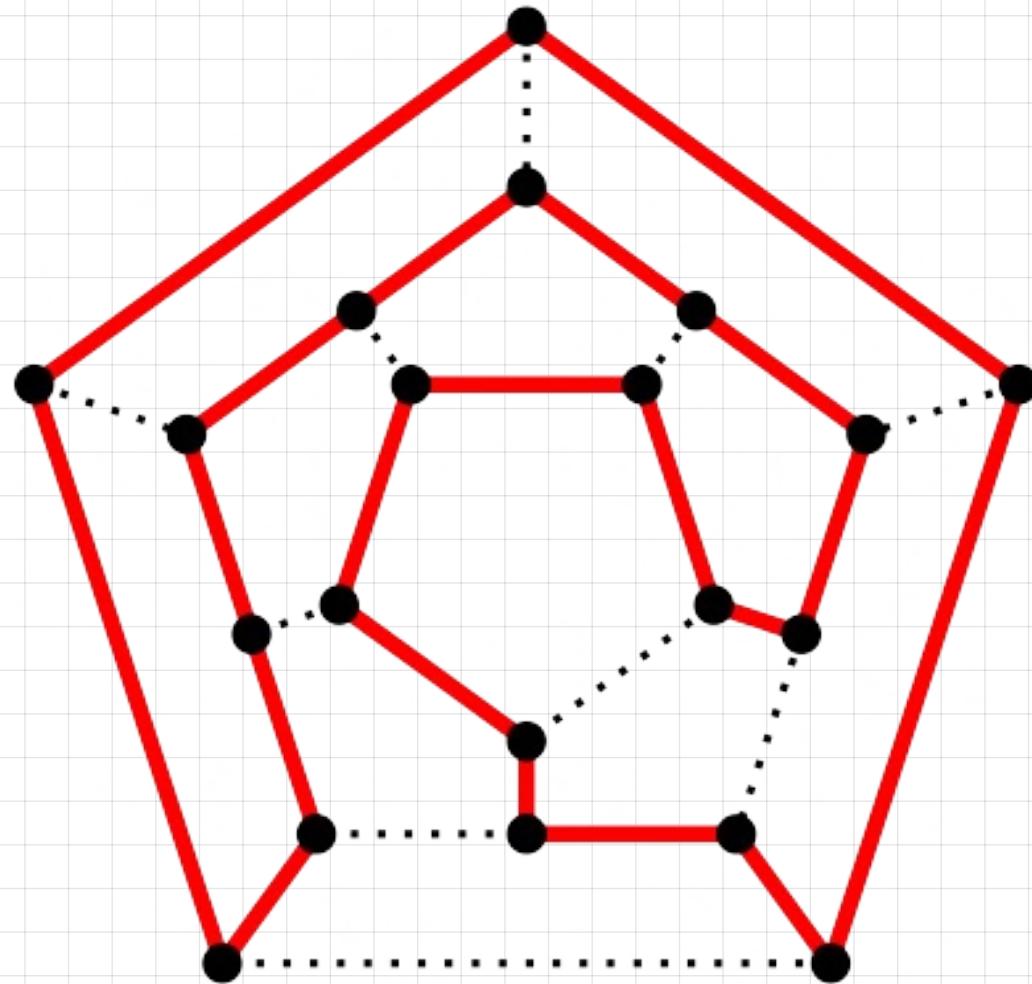
1. Exercise : modify proof of Euler's Thm. to prove :

Connected G has an Eulerian walk starting at s & ending at $t \iff$ all vertices of G except s & t have even degree

2. A Hamilton cycle in G is a cycle in G that visits every vertex exactly once

- no simple characterization

- NP-complete (related to Traveling Salesman)



Example of a graph with a Hamilton cycle
("dodecahedron" graph)

Graph Theory : Topics

- Eulerian tours
- Trees & complete graphs
- Planar graphs
- Connectivity & hypercubes

} next lecture

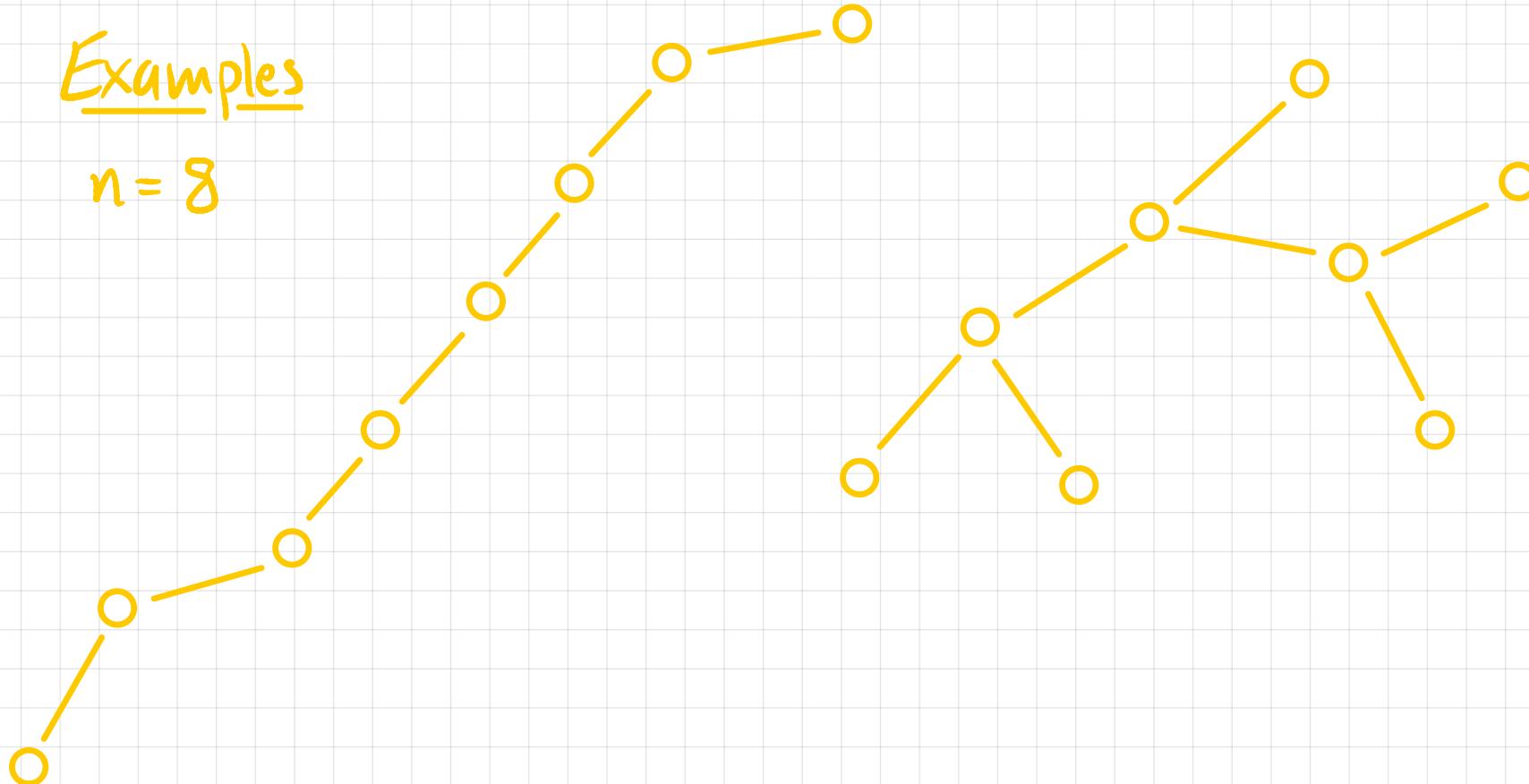
Trees

An undirected graph G with n vertices is a tree if either of the following equiv. conditions holds:

- (i) G is connected & has $n-1$ edges
- (ii) G is connected & has no cycles

Examples

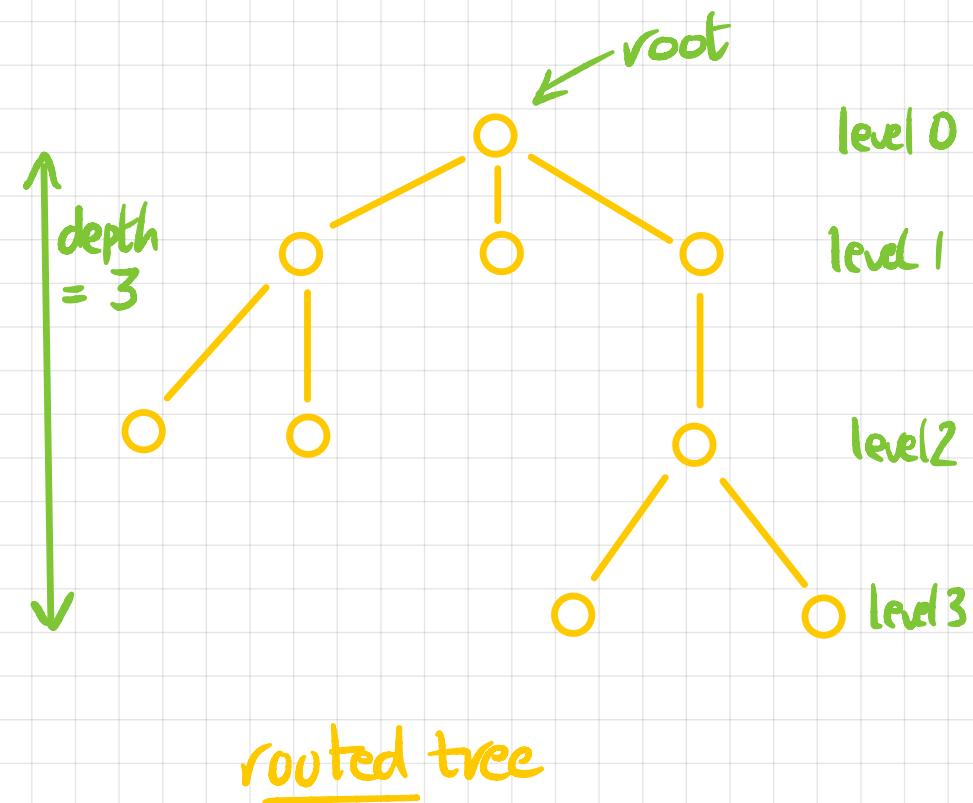
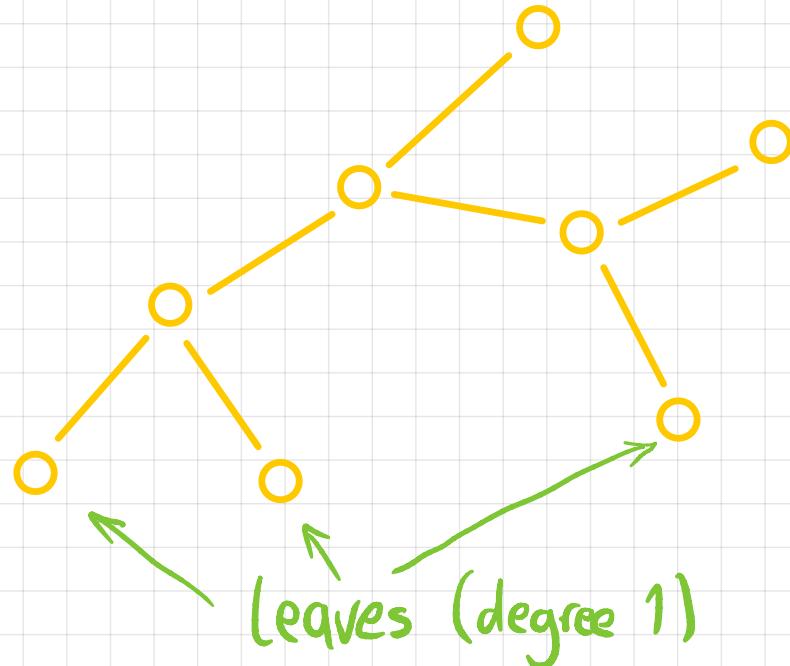
$$n=8$$



Trees show up in :

- data structures (BST, Red-Black trees, ...)
- phylogenetic trees (evolutionary biology, genealogy)
- epidemic models
- decision trees

:



Theorem : G is connected & has $n-1$ edges \iff G is connected & has no cycles

Proof :

Theorem : G is connected & has $n-1$ edges \iff G is connected & has no cycles

Proof :

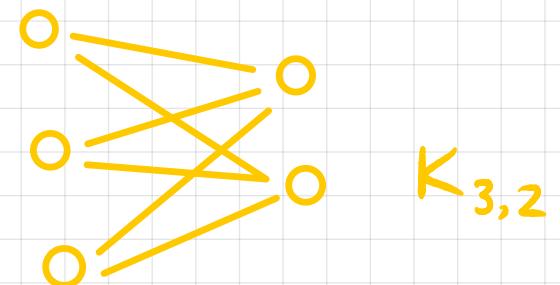
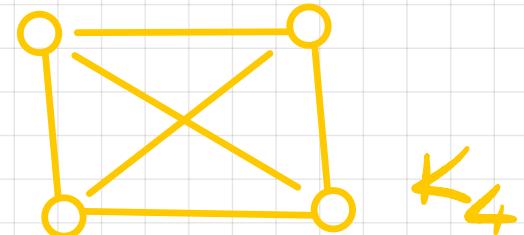
The Complete Graph

The complete graph on n vertices, K_n , is the graph that contains all possible edges (so # of edges is $\frac{n(n-1)}{2}$)

Notes:

1. K_n is unique but \exists many (n^{n-2}) trees on n vertices
2. K_n is maximally connected (need to remove at least $n-1$ edges to disconnect); trees are minimally connected (removing any edge disconnects)
3. Complete bipartite graph $K_{n,m}$:

of edges =



Summary

- Graphs: directed & undirected
- Paths, cycles, walks, tours
- Eulerian tours
- Trees, complete graphs

Next lecture

- Planar graphs
- Hypercubes & connectivity