Midterm 1

7:00-9:00pm, 7 October

Your First Name:	Your Last Name:	
SIGN Your Name:	Your SID Number:	
Your Exam Room:		
Name of Person Sitting on Your Left:		
Name of Person Sitting on Your Right:		
Name of Person Sitting in Front of You:		
Name of Person Sitting Behind You:		

Instructions:

- (a) As soon as the exam starts, please write your student ID in the space provided at the top of every page! (We will remove the staple when scanning your exam.)
- (b) There are **6 double-sided** sheets (12 numbered pages) on the exam. Notify a proctor immediately if a sheet is missing.
- (c) We will not grade anything outside of the space provided for a question (i.e., either a designated box if it is provided, or otherwise the white space immediately below the question). Be sure to write your full answer in the box or space provided! Scratch paper is provided on request; however, please bear in mind that nothing you write on scratch paper will be graded!
- (d) The questions vary in difficulty, so if you get stuck on any question it may help to leave it and return to it later.
- (e) On questions 1-2, you need only give the answer in the format requested (e.g., True/False, an expression, a statement.) An expression may simply be a number or an expression with a relevant variable in it. For short answer questions, correct, clearly identified answers will receive full credit with no justification. Incorrect answers may receive partial credit.
- (f) On questions 3-6, you should give arguments, proofs or clear descriptions if requested. If there is a box you must use it for your answer: answers written outside the box may not be graded!
- (g) You may consult one two-sided "cheat sheet" of notes. Apart from that, you may not look at any other materials. Calculators, phones, computers, and other electronic devices are NOT permitted.
- (h) You may use, without proof, theorems and lemmas that were proved in the notes and/or in lecture.
- (i) You have 120 minutes: there are 6 questions on this exam worth a total of 140 points.

1. True/False [No justification; answer by shading the correct bubble. 2 points per answer; total of 36 points. No penalty for incorrect answers.]

(a) Which of the following is a valid logical equivalence, for arbitrary propositions P, Q, R? Answer **YES** (for valid) or **NO** (for not valid) by shading the appropriate bubble.

YES NO

$$(P \Longrightarrow \neg Q) \equiv \neg (P \land Q)$$
 2pts

$$\bigcirc \quad \bigcirc \quad ((P \land Q) \Longrightarrow R) \equiv (Q \Longrightarrow (\neg P \lor R))$$
 2pts

$$\bigcirc \quad \bigcirc \quad ((P \Longrightarrow R) \land (Q \Longrightarrow R)) \equiv ((P \land Q) \Longrightarrow R)$$
 2pts

(b) Indicate which of the following quantified statements is **TRUE** or **FALSE** by shading the appropriate bubble.

TRUE FALSE

$$\bigcirc \qquad (\forall a, b \in \mathbb{N})(\exists c \in \mathbb{N})((a < b) \Rightarrow (a < c < b)).$$
 2pts

$$\bigcirc \qquad \bigcirc \qquad \neg (\exists a, b \in \mathbb{N})(\sqrt{2} = \frac{a}{b}).$$
 2pts

$$\bigcirc \qquad (\forall n \in \mathbb{N})(\exists m \in \mathbb{N})((m > n) \land \neg(\exists a, b \in \mathbb{N})((a > 1) \land (b > 1) \land (m = ab))).$$
 2pts

$$\bigcirc \qquad \bigcirc \qquad (\forall x, y, z \in \mathbb{Z})((x \le y) \Rightarrow (xz \le yz)).$$
 2pts

(c) Which of the following are valid proof strategies? Answer **YES** (for valid) or **NO** (for not valid) by shading the appropriate bubble.

YES NO

- \bigcap P(0), P(1) and P(2) hold, and $(\forall k \geq 2)(P(k) \Rightarrow P(k+2))$ holds. Hence P(n) holds for all 2pts $n \in \mathbb{N}$.
- \bigcirc P(0) holds and for all $k \in \mathbb{N}$ the following two implications hold: $P(k) \Rightarrow P(2k)$; $P(2k) \Rightarrow 2pts$ P(2k+1). Hence P(n) holds for all $n \in \mathbb{N}$.
- \bigcirc P(0) holds and for all $k \in \mathbb{N}$ the following three implications hold: $P(2k) \Rightarrow P(2k+1)$; 2pts $P(-2k+1) \Rightarrow P(-2k)$; and $P(k) \Leftrightarrow P(-k)$. Hence P(n) holds for all $n \in \mathbb{Z}$.

(d) Answer each of the following questions **TRUE** or **FALSE** by shading the appropriate bubble.

TRUE FALSE

\bigcirc	0	In a stable marriage instance, if a woman W is the least favorite in every man's preference list, then changing W 's preference list cannot change the pairing output by the traditional Stable Marriage Algorithm.	2pts
0	0	In a stable marriage instance, if a man M is the least favorite in every woman's preference list, then changing M 's preference list cannot change the pairing output by the traditional Stable Marriage Algorithm.	2pts
\bigcirc	\bigcirc	Every planar drawing of a graph G has the same number of faces.	2pts
\bigcirc	\bigcirc	If a graph has 10 vertices and 25 edges, then it cannot be planar.	2pts
\bigcirc	\bigcirc	Every tree with at least four vertices must have at least three leaves.	2pts
\bigcirc	\bigcirc	Any two longest (simple) paths in a connected, undirected graph must have at least one vertex in common.	2pts
\bigcirc	\bigcirc	The equation $5x \equiv 15 \pmod{93}$ has a unique solution for $x \pmod{93}$.	2pts
\bigcirc	\bigcirc	The equation $3x \equiv 18 \pmod{93}$ has a unique solution for x modulo 93.	2pts

(f) Solve the equation $4x \equiv 7 \pmod{11}$ for x. Your answer should be an integer in $\{0, 1, \dots, 10\}$.

3pts

- (g) Compute $10^{11^{12^{13}}}$ (mod 11). Your answer should be an integer in $\{0,1,\ldots,10\}$. 3pts

 (h) Compute $1^{16}+2^{16}+\cdots+99^{16}+100^{16}$ (mod 17). Your answer should be an integer in $\{0,1,\ldots,16\}$. 3pts

 (i) Let n be a positive integer. What is $\gcd(4n^2+1,2n+1)$? 3pts

 (j) Find $\gcd(476,153)$. 3pts
- (k) Find the inverse of 24 mod 53. Your answer should be an integer in $\{0, 1, \dots, 52\}$.
- (1) Bob sets up an RSA scheme based on primes p=5, q=11 and exponent e=3. Alice wants to send 3pts Bob the message 20 (viewed as an integer mod 55). What is the encrypted message that Alice sends to Bob?

- **3. Proofs** [All parts to be justified. Total of 20 points.]
 - (a) Let a be a natural number. Prove that if $a^2 3a + 1$ is even then a is odd.

4pts

(b) Prove that $\sqrt{ab} \le \frac{a+b}{2}$ for all positive real numbers a,b.

4pts

(c) Prove that there do not exist integers m, n such that $m^2 + 4n + 1 = 0$.

5pts

(d) Prove by induction that, for all natural numbers n, $2^{3n+1} + 5$ is divisible by 7.

7pts

4. Stable Marriage [All parts to be justified unless otherwise stated. Total of 18 points.] Consider the following preference lists for three men 1, 2, 3, and three women A, B, C.

Man	Women			Woman	M	Men	
1	A	В	С	A	3		
2	В	A	C	В	1	2	3
3	В	C	A	C	2	3	1

- (a) Find the female-optimal pairing. Write your answer in the boxes provided; no justification is needed. 3pts

 Man 1 is paired with: Man 2 is paired with: Man 3 is paired with: Man 4pts

 (c) There is one other stable pairing, different from both the female-optimal and female-pessimal ones. What is this pairing? Write your answer in the boxes provided; no justification is needed. [Hint: Use parts (a) and (b) to narrow your search.]

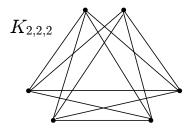
 Man 1 is paired with: Man 2 is paired with: Man 3 is paired with: Man 3 is paired with:
- (d) Suppose now that the women know all the preference lists, and wish to collude in order to achieve 4pts the female-optimal pairing you found in part (a), even though the traditional algorithm (in which the men propose) will be used. Prove that there is no way in which the women can jointly revise their preference lists so as to achieve this pairing. [Hint: Note that the men's lists remain fixed. You should answer this question without considering all possibilities for the womens' lists. You may use without proof any result you have seen in class or in the homeworks, provided you state it clearly.]

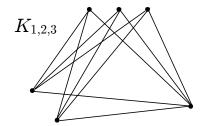
(e) Now suppose that the women are given the additional option of omitting some men from their preference lists altogether (i.e., a woman may rank only some of the men, declaring the others "unacceptable"). Suppose also that we again run the traditional (men-proposing) algorithm, except that a woman rejects all proposals from her unacceptable men (possibly leaving her with no man on a string). Explain how the women can now achieve the female-optimal pairing from part (a), and justify why this works.

3pts

5. Complete Tripartite Graphs. [All parts to be justified unless otherwise stated. Total of 18 points.]

Let ℓ, m, n be positive integers with $\ell \leq m \leq n$. The complete tripartite graph $K_{\ell,m,n}$ consists of three disjoint groups of vertices, L, M, N, with $|L| = \ell$, |M| = m, and |N| = n. There are no edges between any two vertices in the same group, and there is one edge between every pair of vertices belonging to different groups. The figure below shows two examples: $K_{2,2,2}$ and $K_{1,2,3}$.





(a) How many *edges* are there in $K_{\ell,m,n}$? Write your answer as a function of ℓ,m,n in the box provided; *3pts* no justification is needed. [Hint: Check your answer on the two examples above!]



(b) For what values of ℓ, m, n does there exist an Eulerian tour in $K_{\ell,m,n}$? No justification is needed.

(c) Prove that $K_{\ell,m,n}$ does **not** have a Hamiltonian cycle if $n > \ell + m$. [Recall that a Hamiltonian cycle 5pts in a graph G is a cycle that visits every vertex of G exactly once.]

Your SID Number: Page 11 (d) It is known that $K_{2,2,2}$ is planar (even though this is not obvious from the above drawing). How many 2pts faces are there in any planar drawing of $K_{2,2,2}$? Write your answer in the box provided; no justification is needed. (e) For which values of ℓ, m, n is $K_{\ell,m,n}$ planar? Justify your answer. [Note: You may use the fact stated 5pts in part (d) that $K_{2,2,2}$ is planar.]

6. RSA [All parts to be justified unless otherwise stated. Total of 12 pts.]

Alice sends messages to Bob using his public key (N,e), where N=pq is the product of two large primes, and e=3. Eve (as usual) is listening in but is unable to decrypt any of the messages. One day, just for fun, Bob decides to share his private key, d, with Eve. Of course, this means that Eve is now able to decrypt all messages sent to Bob. However, Bob isn't too worried as he plans to change the value e in his public key to some other value e', reasoning that Eve won't be able to compute Bob's corresponding new private key d' because she still doesn't know p and q. In this problem, you are asked to show that Bob is mistaken, because in fact, using her knowledge of d, Eve can factor Bob's public key N=pq and thus do anything Bob can do.

(a) Show that 3d - 1 = k(p - 1)(q - 1) for k = 1 or k = 2.

4pts

(b) How can Eve determine the value of k in part (a)?

4pts

(c) Using parts (a) and (b), show how Eve can factor N efficiently. [Hint: You may want to give an explicit 4pts formula for p (or for q) in terms of N and d. Alternatively, you may write down equations that can be solved for p (or q).]

[End of Exam!]